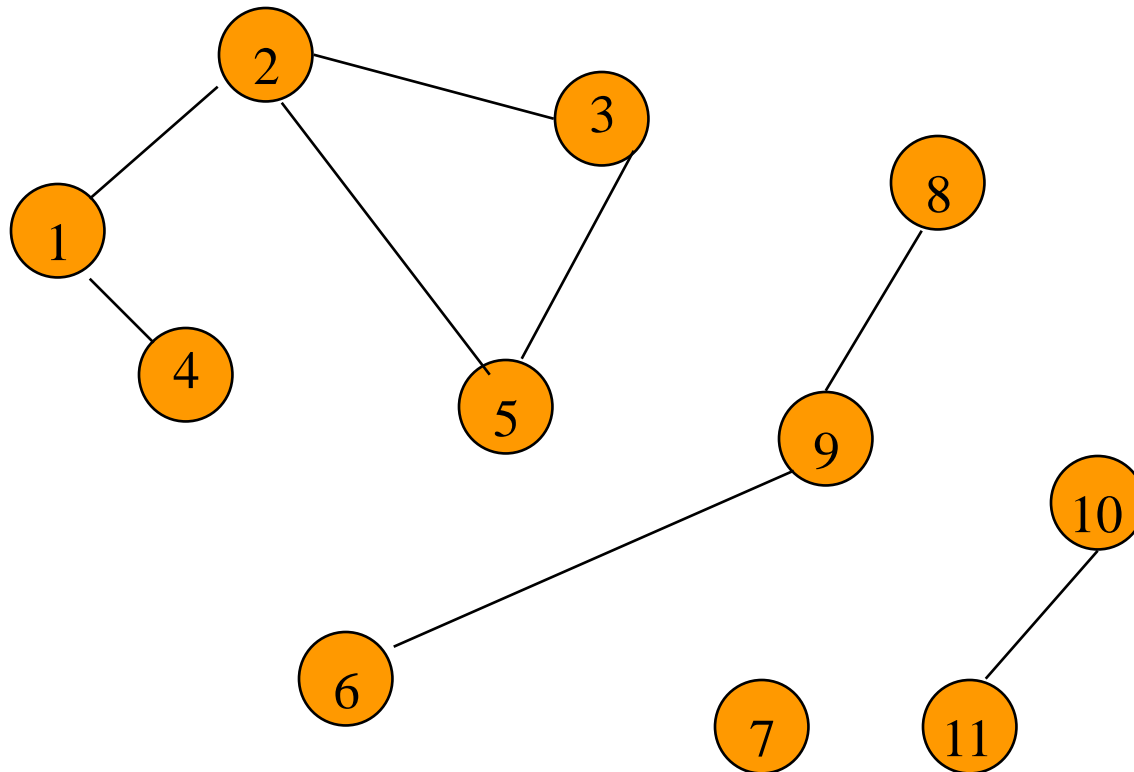


# Graph Search Methods

# Graph Search Methods

- A search method starts at a given vertex **v** and visits/labels/marks every vertex that is reachable from **v**.



# Graph Search Methods

- Many graph problems can be solved using a search method.
  - Path from one vertex to another.
  - Is the graph connected?
  - Find a spanning tree.
  - ...
- Commonly used search methods:
  - Breadth-first search.
  - Depth-first search.

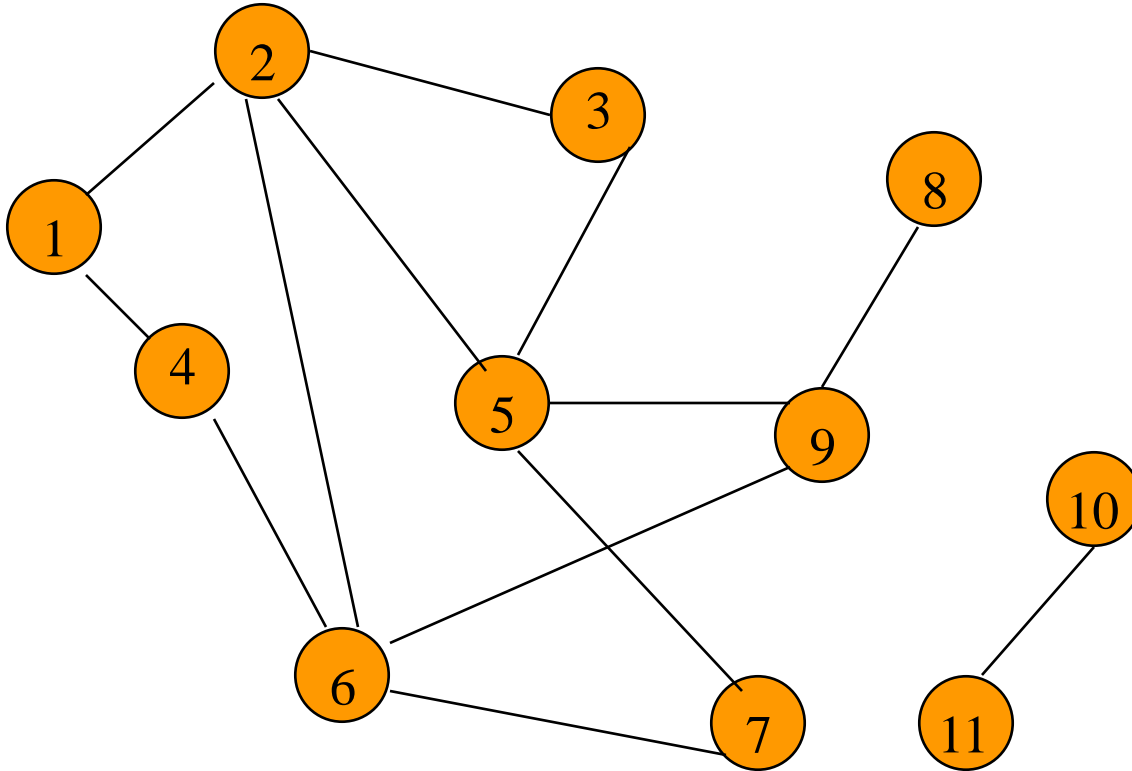
# Graph Search Methods

- In both DFS and BFS, the nodes of the undirected graph are visited in a systematic manner so that every node is visited exactly one.
- Both BFS and DFS give rise to a tree:
  - When a node  $x$  is visited, it is labeled as visited, and it is added to the tree
  - If the traversal got to node  $x$  from node  $y$ ,  $y$  is viewed as the parent of  $x$ , and  $x$  a child of  $y$

# Breadth-First Search

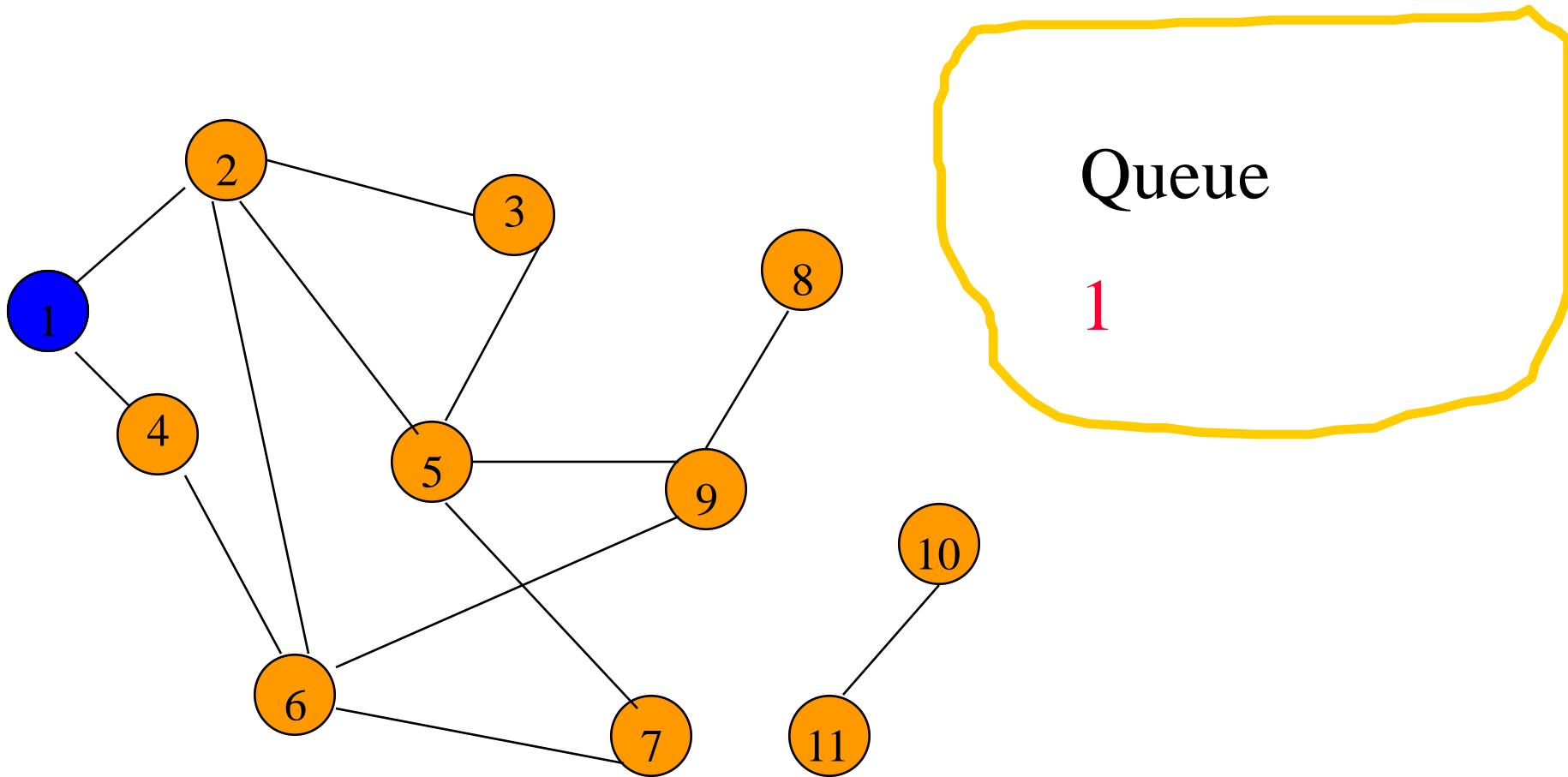
- Visit start vertex and put into a queue.
- Repeatedly remove a vertex from the queue, visit its unvisited adjacent vertices, put newly visited vertices into the queue until the queue is empty.

# Breadth-First Search Example



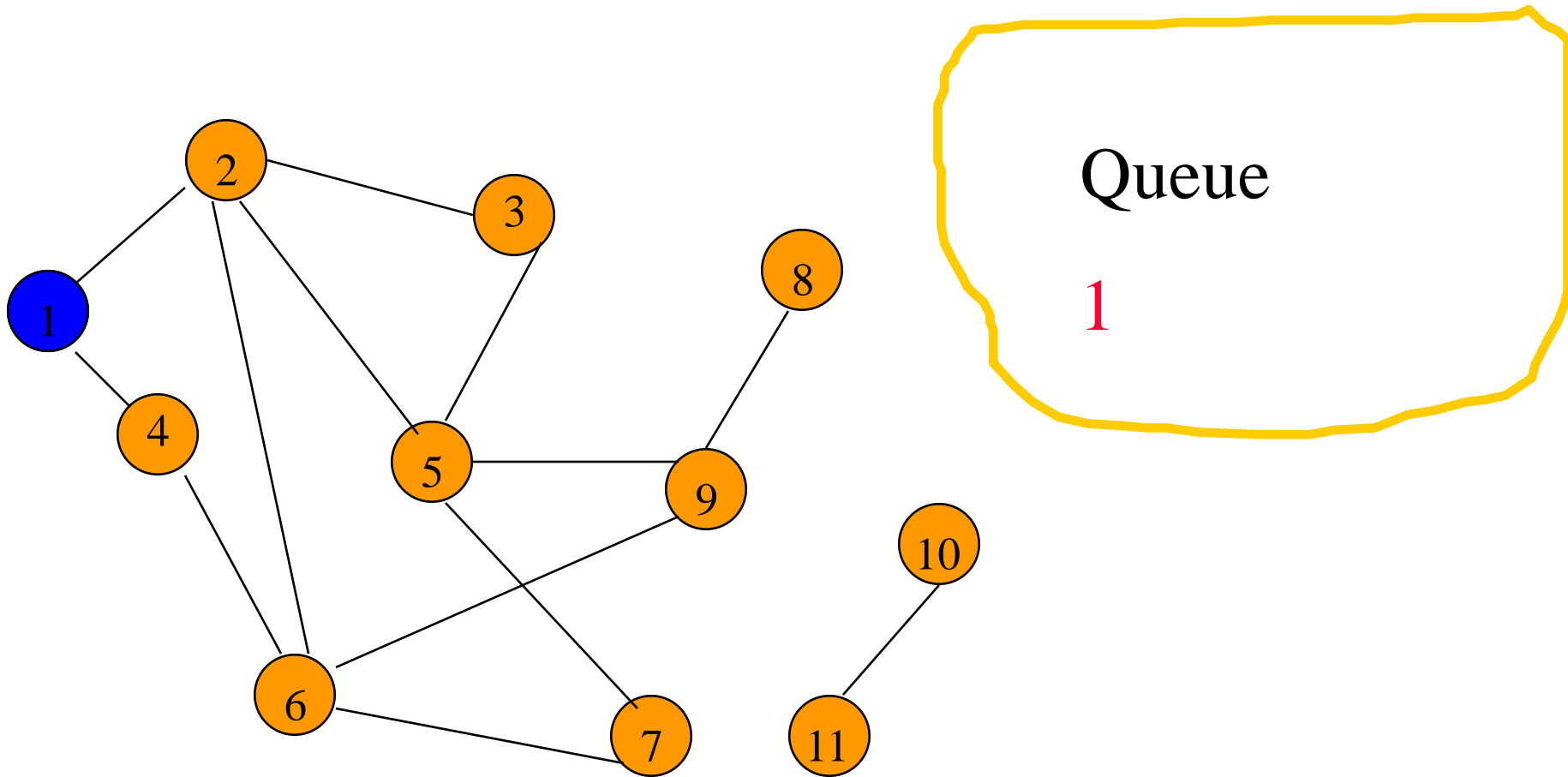
Start search at vertex **1**.

# Breadth-First Search Example



Visit/mark/label start vertex and put in a queue.

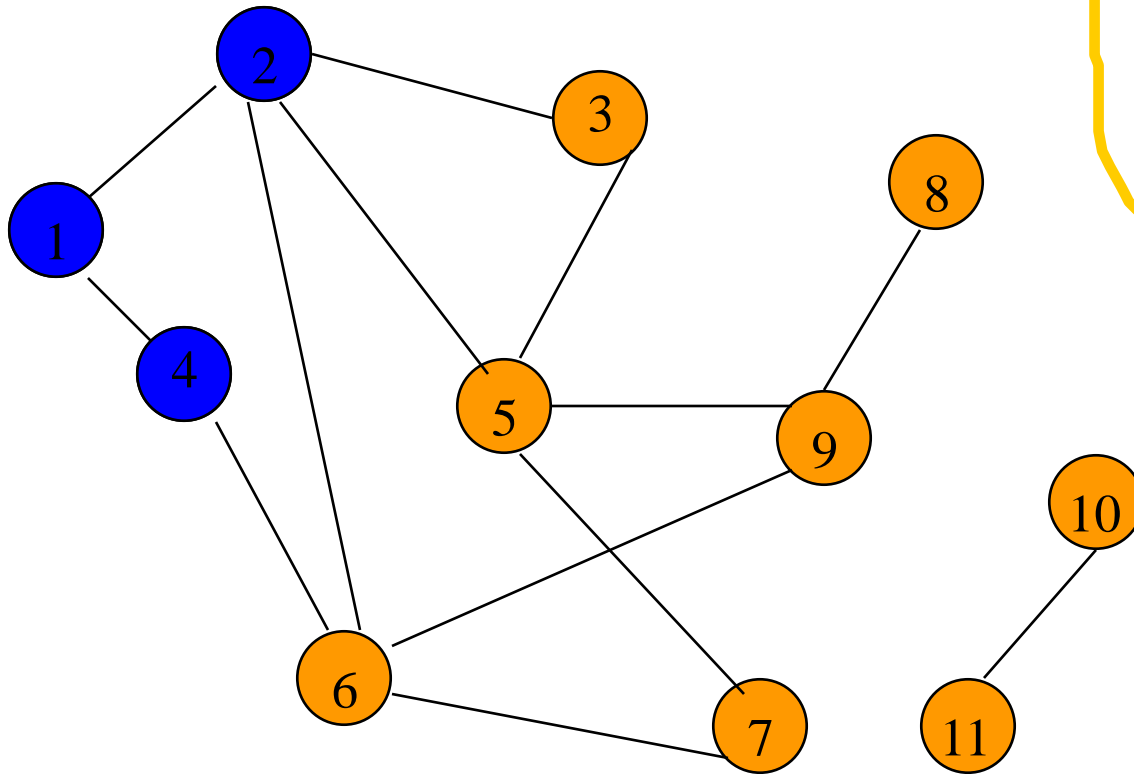
# Breadth-First Search Example



Remove 1 from  $Q$ ; visit adjacent unvisited vertices;  
put in  $Q$ .



# Breadth-First Search Example

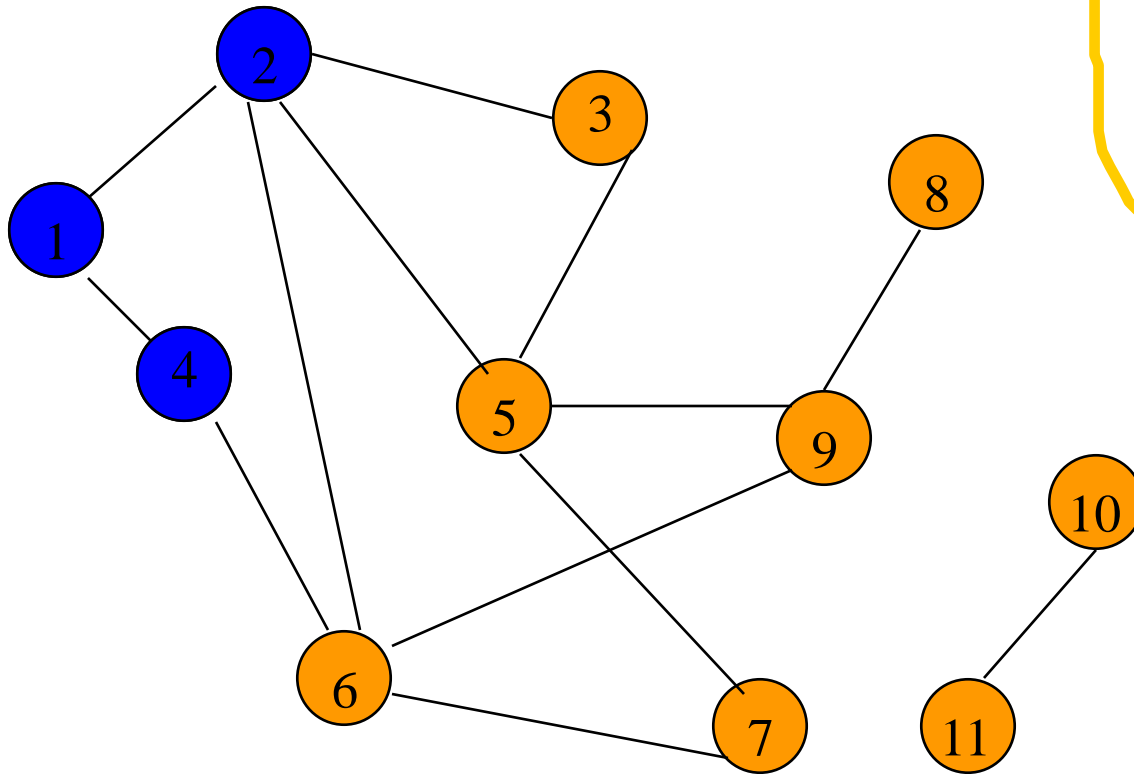


Queue

2 4

Remove 1 from Q; visit adjacent unvisited vertices;  
put in Q.

# Breadth-First Search Example

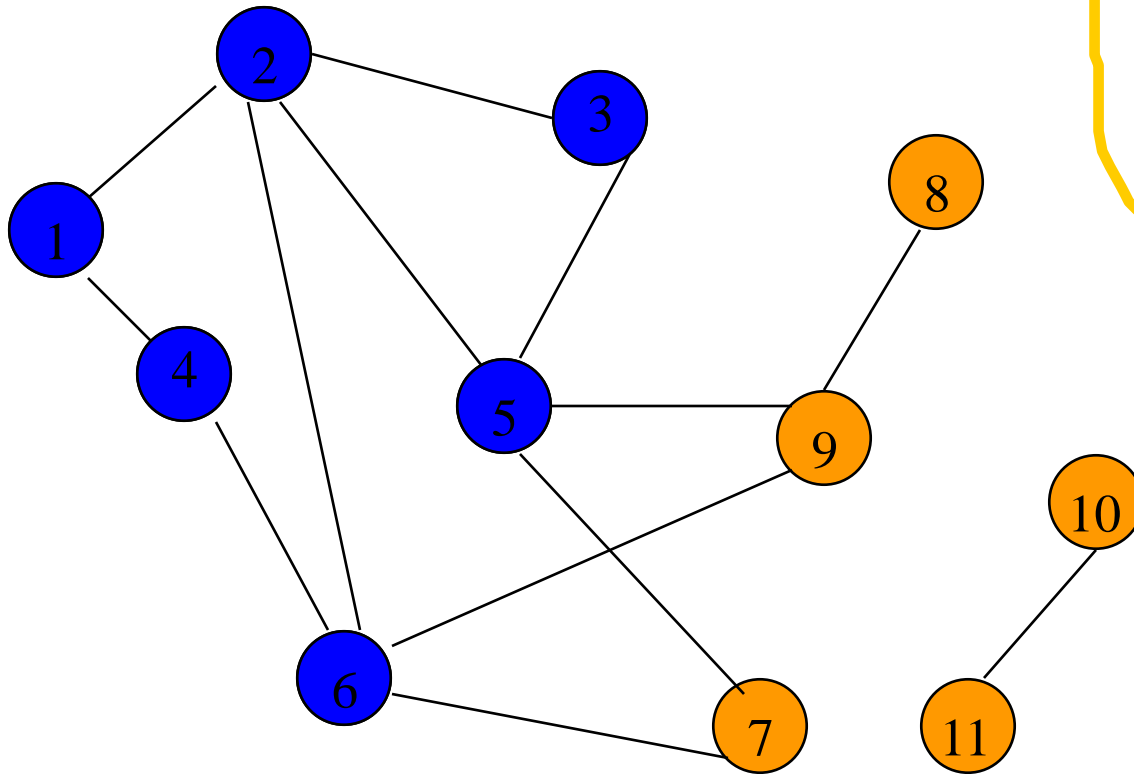


Queue

2 4

Remove 2 from Q; visit adjacent unvisited vertices;  
put in Q.

# Breadth-First Search Example

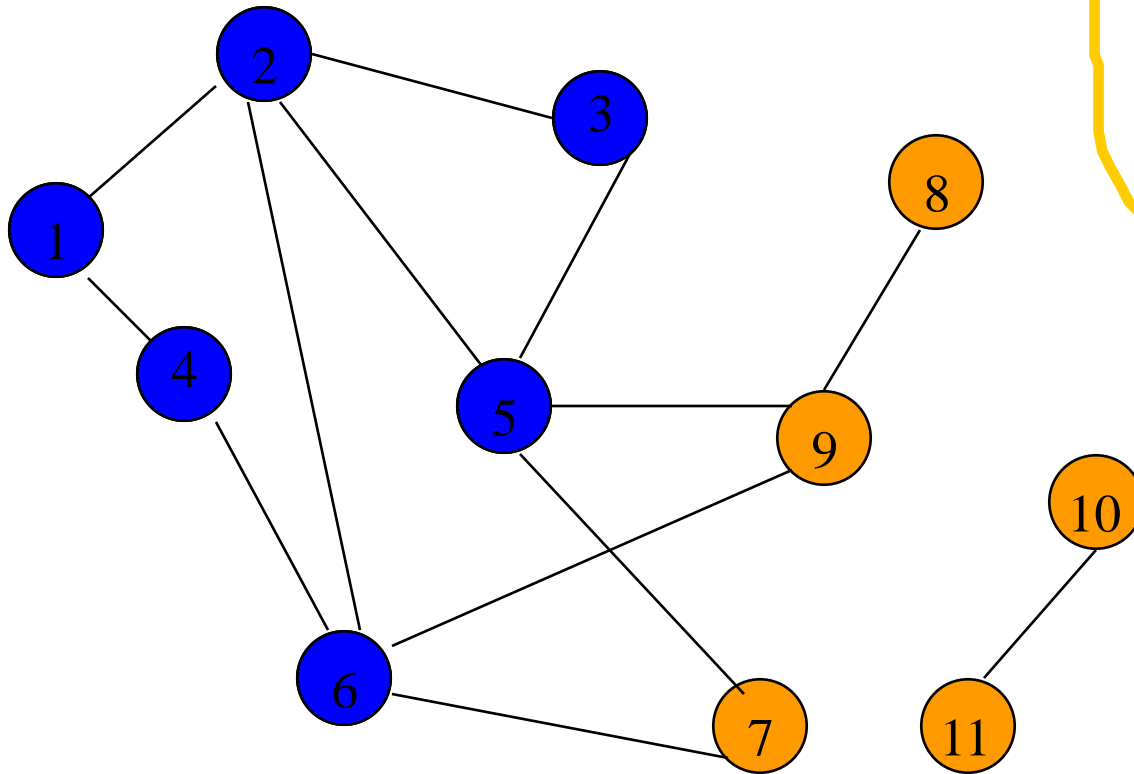


Queue

4 5 3 6

Remove 2 from Q; visit adjacent unvisited vertices;  
put in Q.

# Breadth-First Search Example

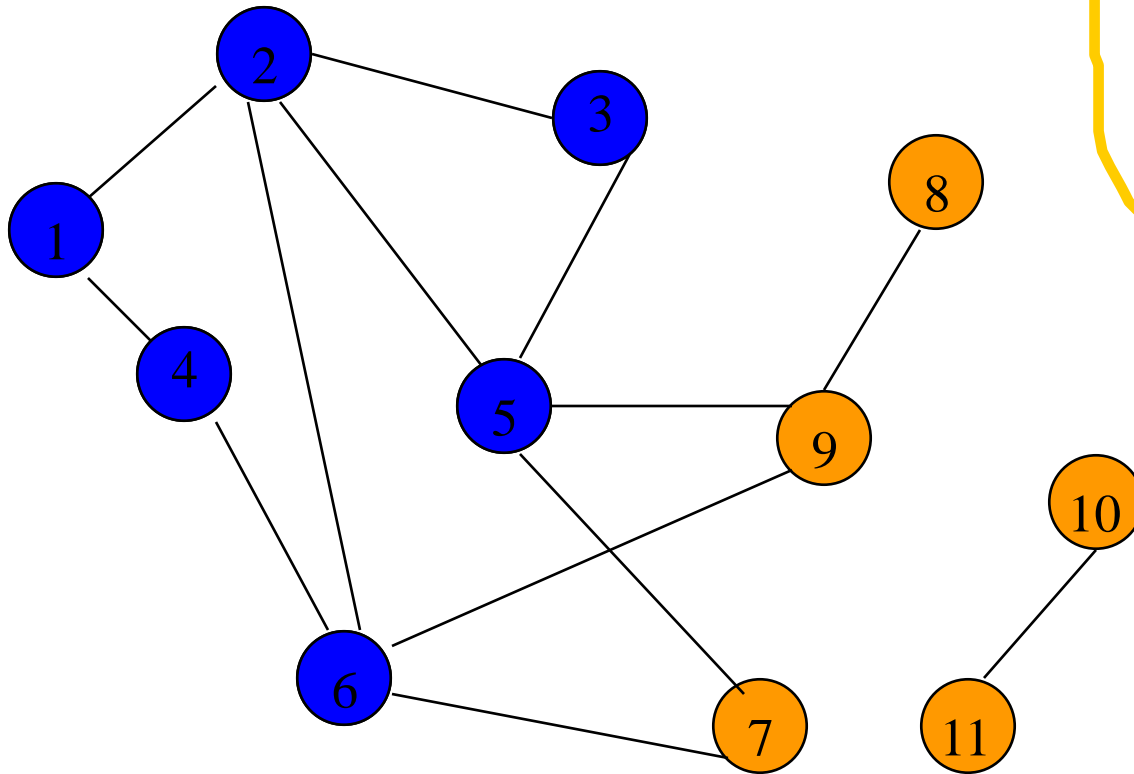


Queue

4 5 3 6

Remove 4 from **Q**; visit adjacent unvisited vertices;  
put in **Q**.

# Breadth-First Search Example

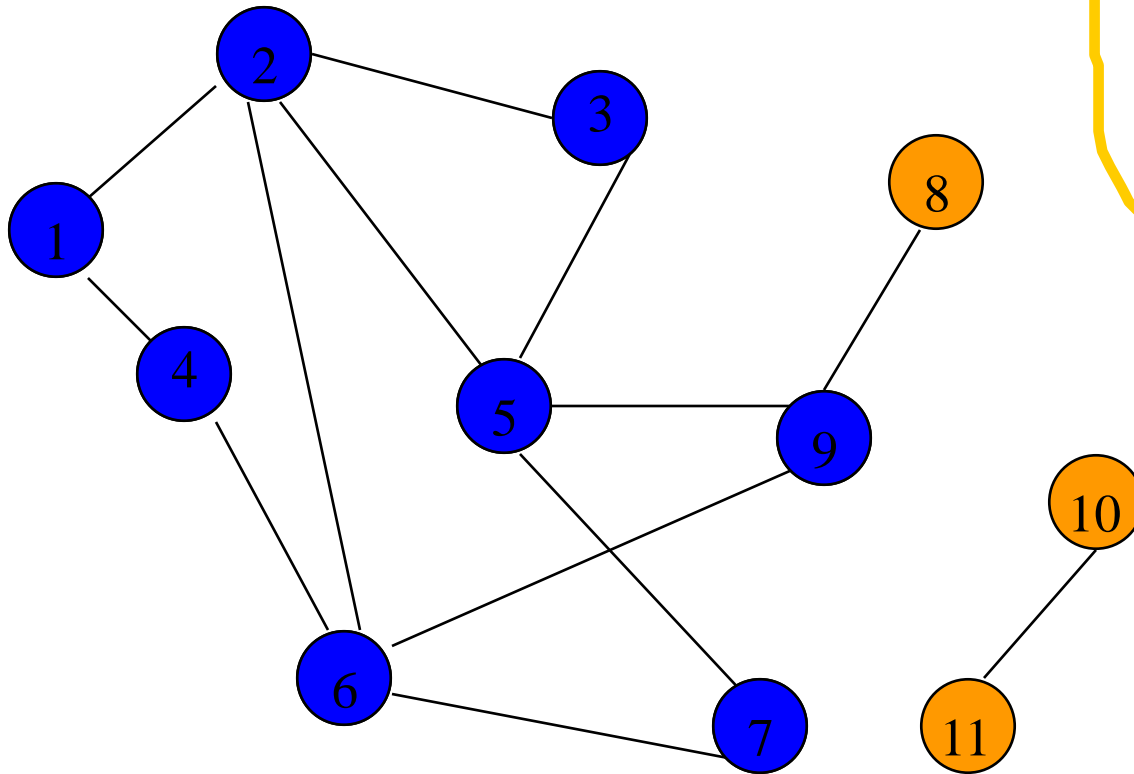


Queue

5 3 6

Remove **5** from **Q**; visit adjacent unvisited vertices;  
put in **Q**.

# Breadth-First Search Example

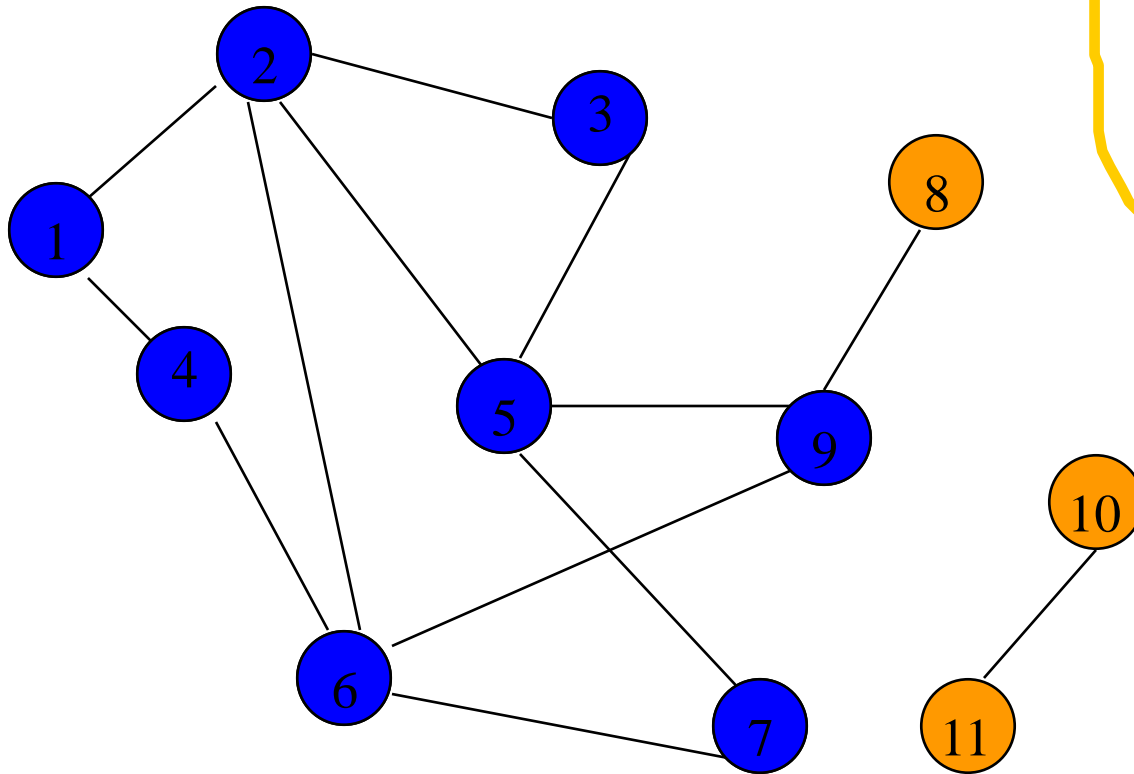


Queue

3 6 9 7

Remove **5** from **Q**; visit adjacent unvisited vertices;  
put in **Q**.

# Breadth-First Search Example

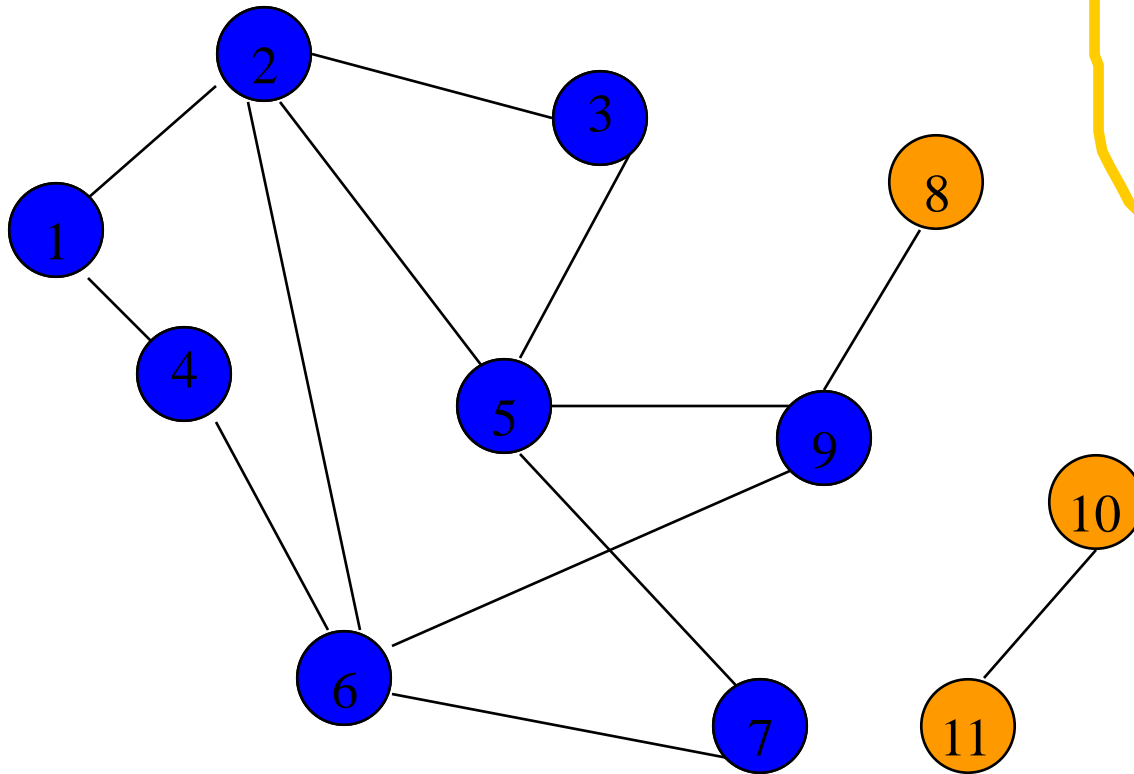


Queue

3 6 9 7

Remove 3 from **Q**; visit adjacent unvisited vertices;  
put in **Q**.

# Breadth-First Search Example



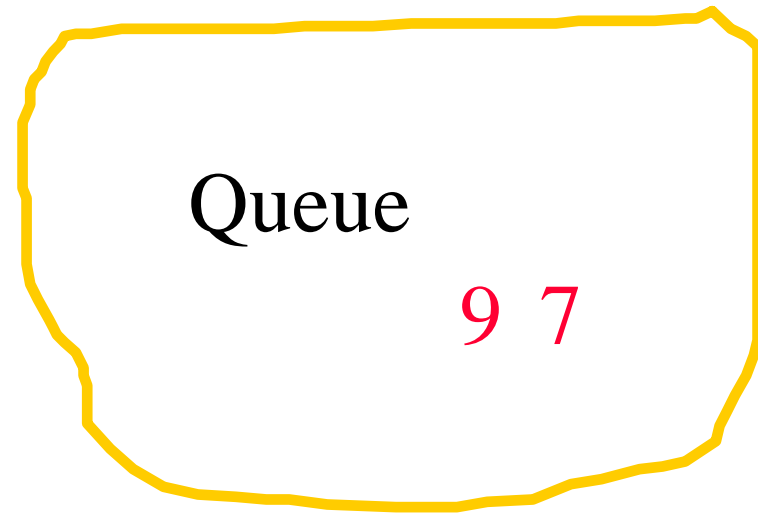
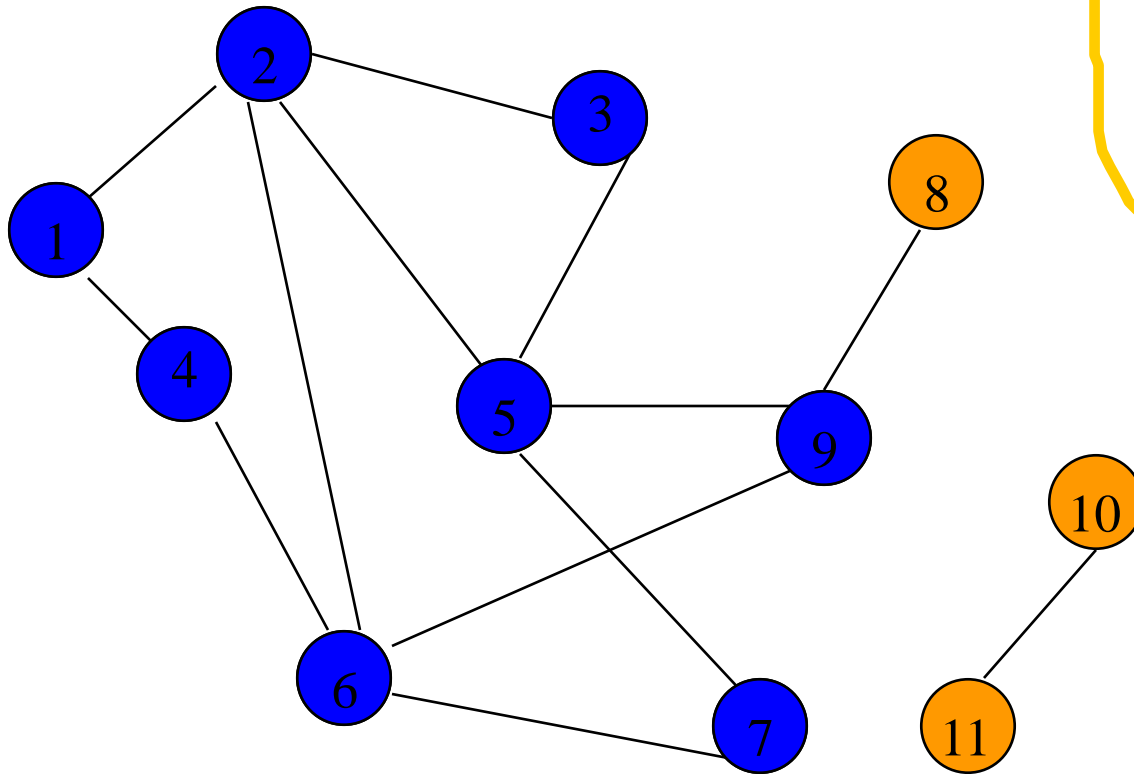
Queue

6 9 7

Remove 6 from Q; visit adjacent unvisited vertices;  
put in Q.

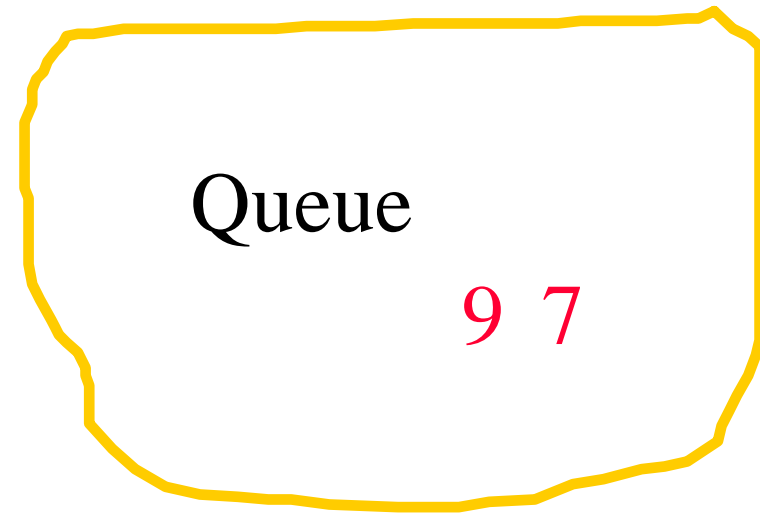
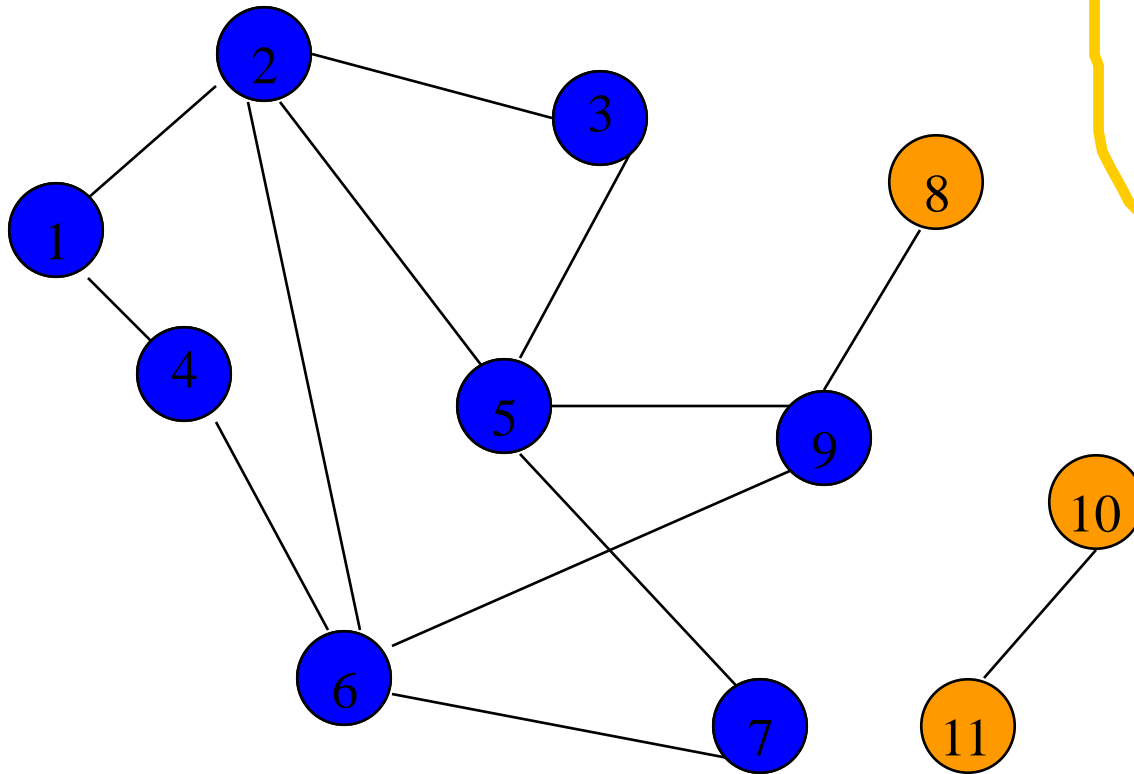


# Breadth-First Search Example



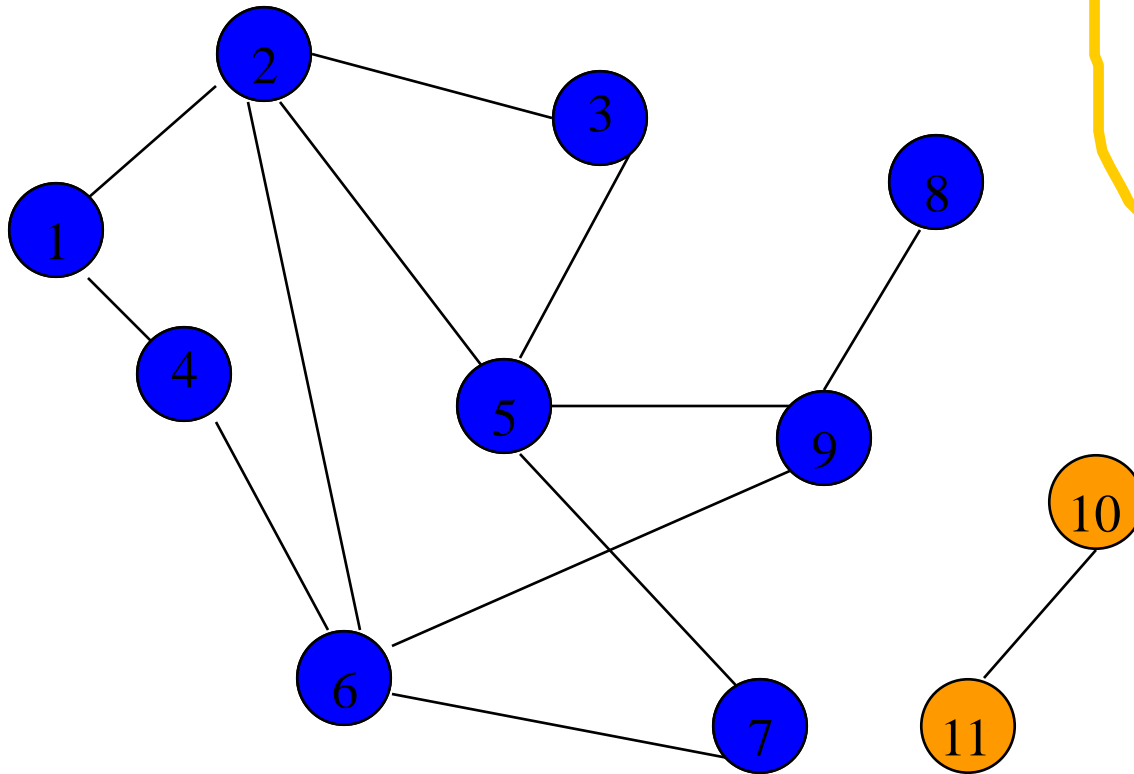
Remove 6 from Q; visit adjacent unvisited vertices;  
put in Q.

# Breadth-First Search Example



Remove 9 from Q; visit adjacent unvisited vertices;  
put in Q.

# Breadth-First Search Example

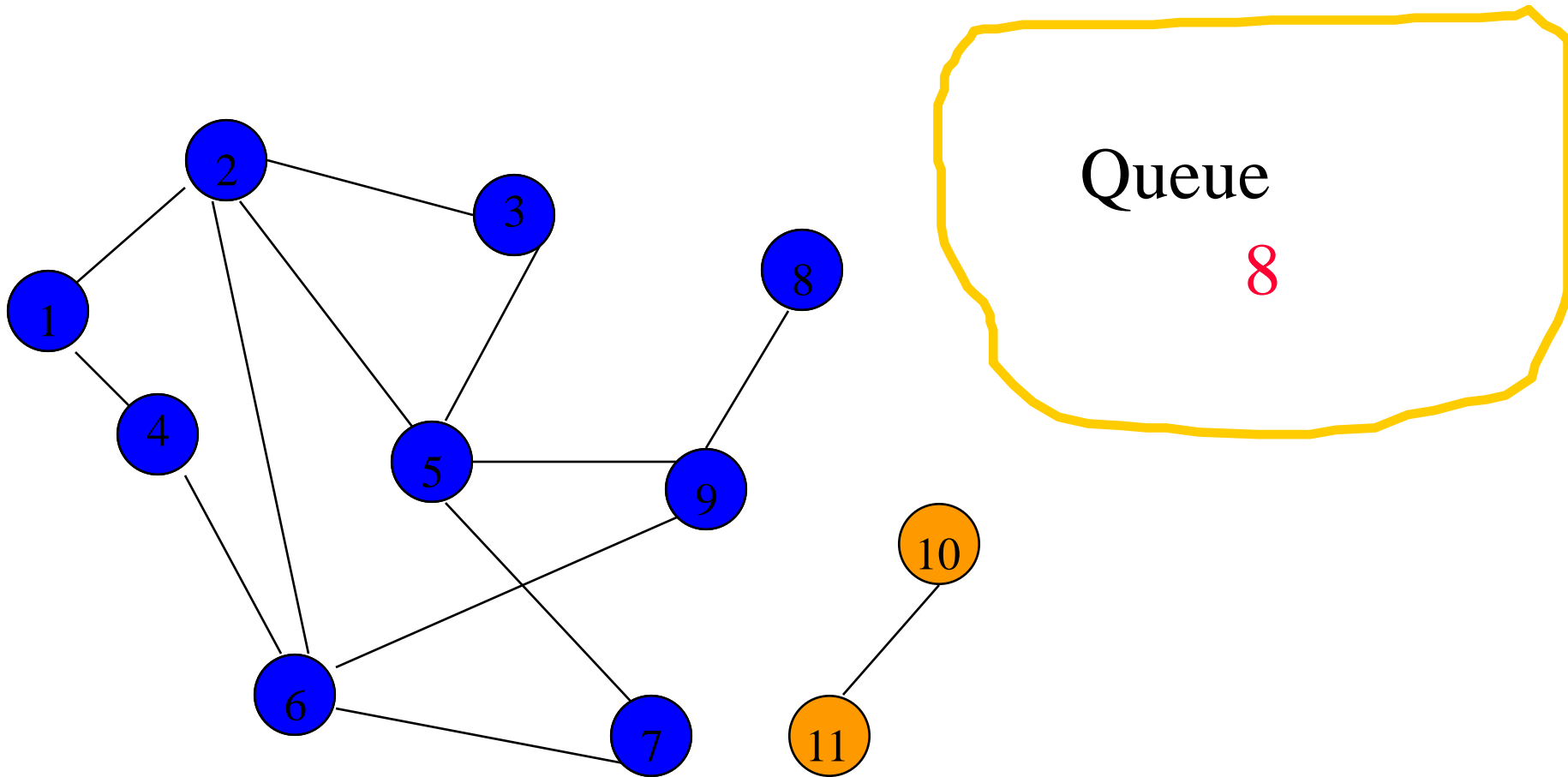


Queue

7 8

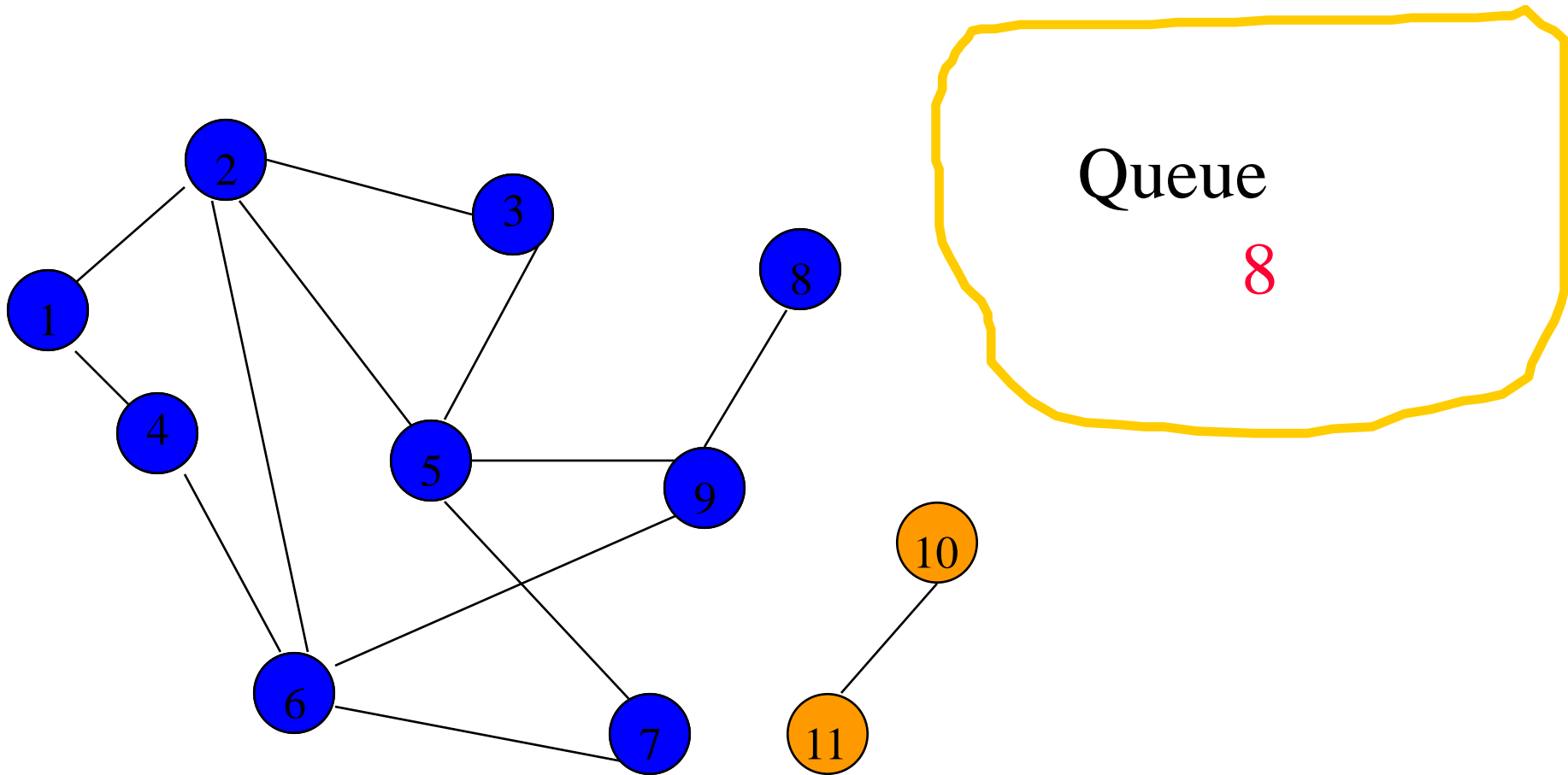
Remove 9 from Q; visit adjacent unvisited vertices;  
put in Q.

# Breadth-First Search Example



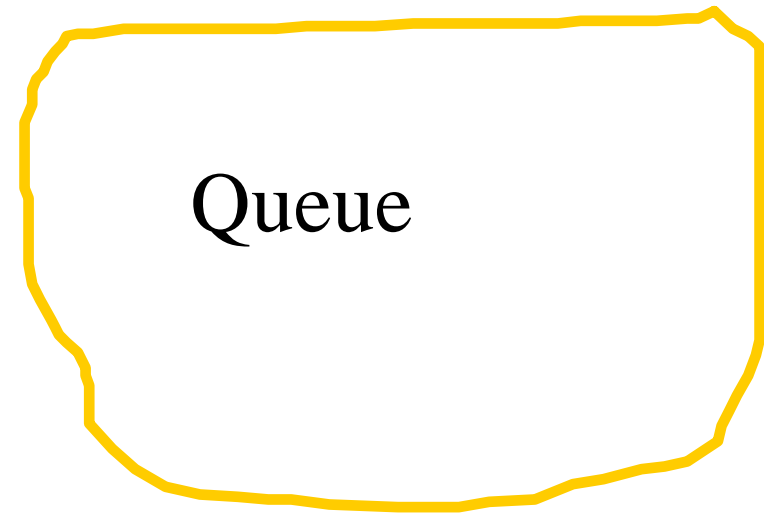
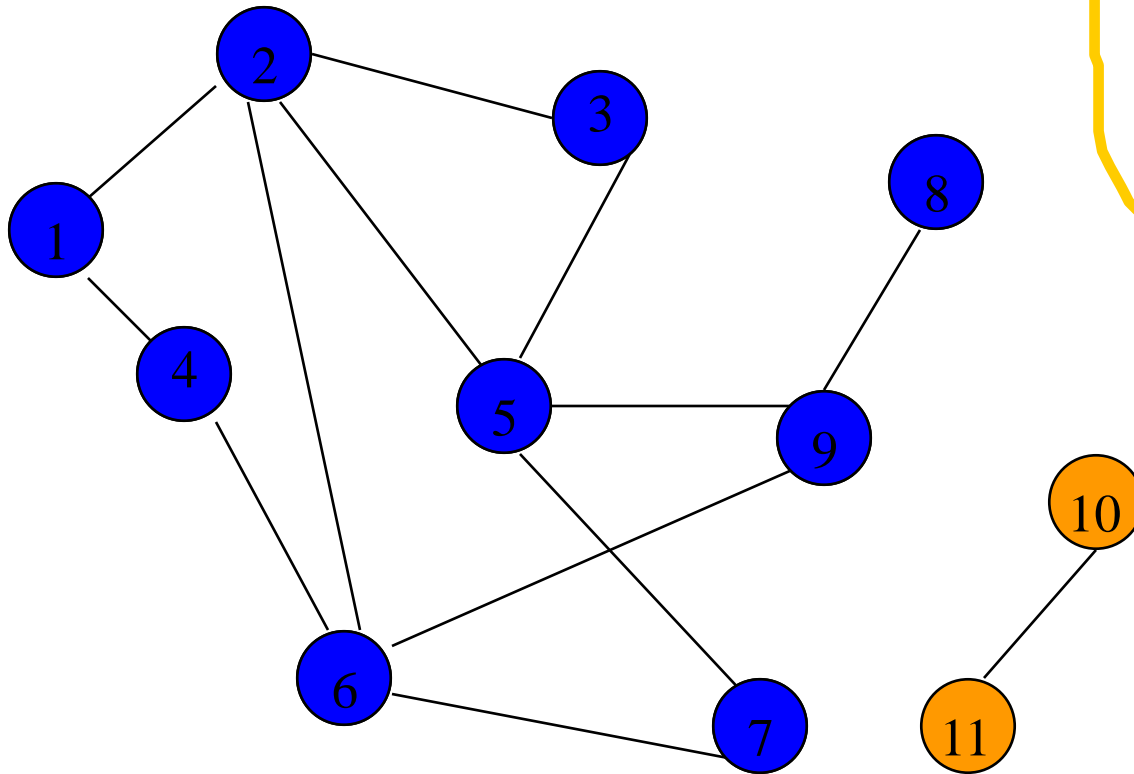
Remove **7** from **Q**; visit adjacent unvisited vertices;  
put in **Q**.

# Breadth-First Search Example



Remove 8 from  $Q$ ; visit adjacent unvisited vertices;  
put in  $Q$ .

# Breadth-First Search Example



Queue is empty. Search terminates.

# Breadth-First Search Property

- All vertices reachable from the start vertex (including the start vertex) are visited.

# Time Complexity



- Each visited vertex is put on (and removed from) the queue exactly once.
- When a vertex is removed from the queue, we examine its adjacent vertices.
  - $O(\text{vertex degree})$  if adjacency lists are used
- Total time, when adjacency lists are used:
  - $O(n + \text{sum of degrees of the vertices in the component})$
  - $= O(n + \text{number of edges in the component})$
  - $= O(n + m)$  if the graph is connected.



# Depth-First Search

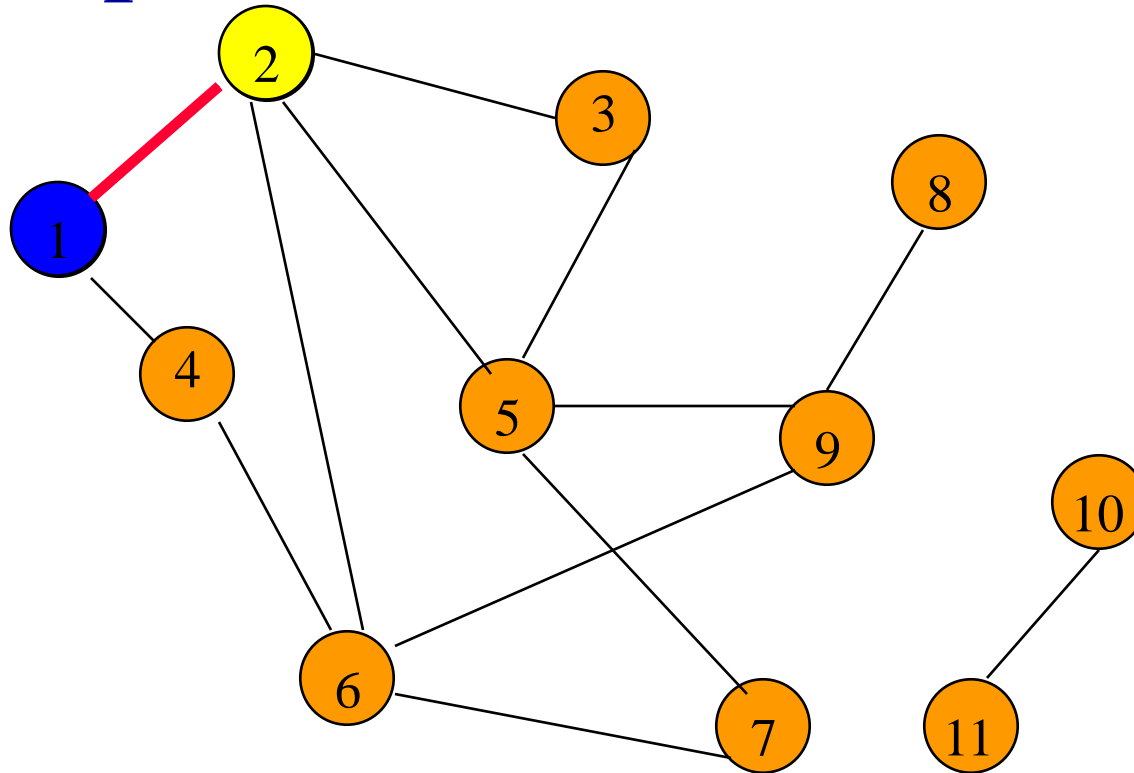
- DFS follows the following rules:
  1. Select an unvisited node  $x$ , visit it, and treat as the **current node**
  2. Find an unvisited neighbor of the current node, visit it, and make it the new current node;
  3. If the current node has no unvisited neighbors, **backtrack** to the its parent, and make that parent the new current node;
  4. Repeat steps 3 and 4 until no more nodes can be visited.
  5. If there are still unvisited nodes, repeat from step 1.

# Depth-First Search

pseudo code

```
depthFirstSearch(v)  
{  
    Label vertex v as reached.  
    for (each unreached vertex u adjacent to v)  
        depthFirstSearch(u);  
}
```

# Depth-First Search Example

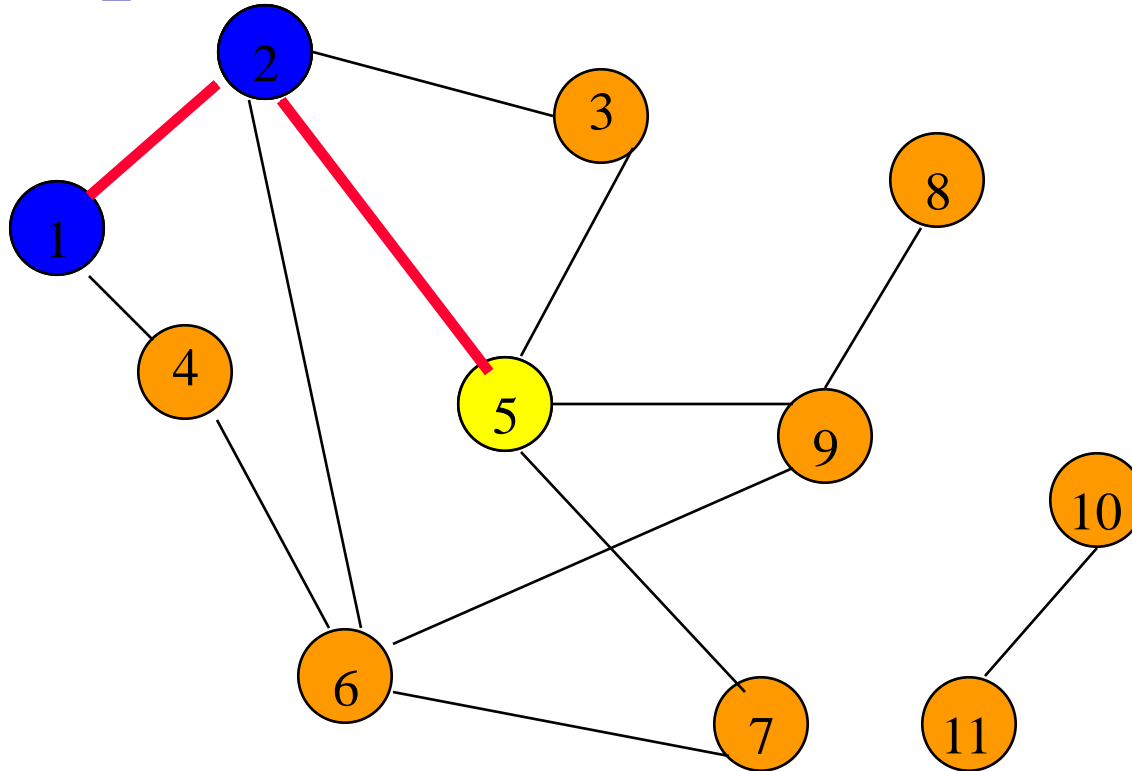


Start search at vertex **1**.

Label vertex **1** and do a depth first search from either **2** or **4**.

Suppose that vertex **2** is selected.

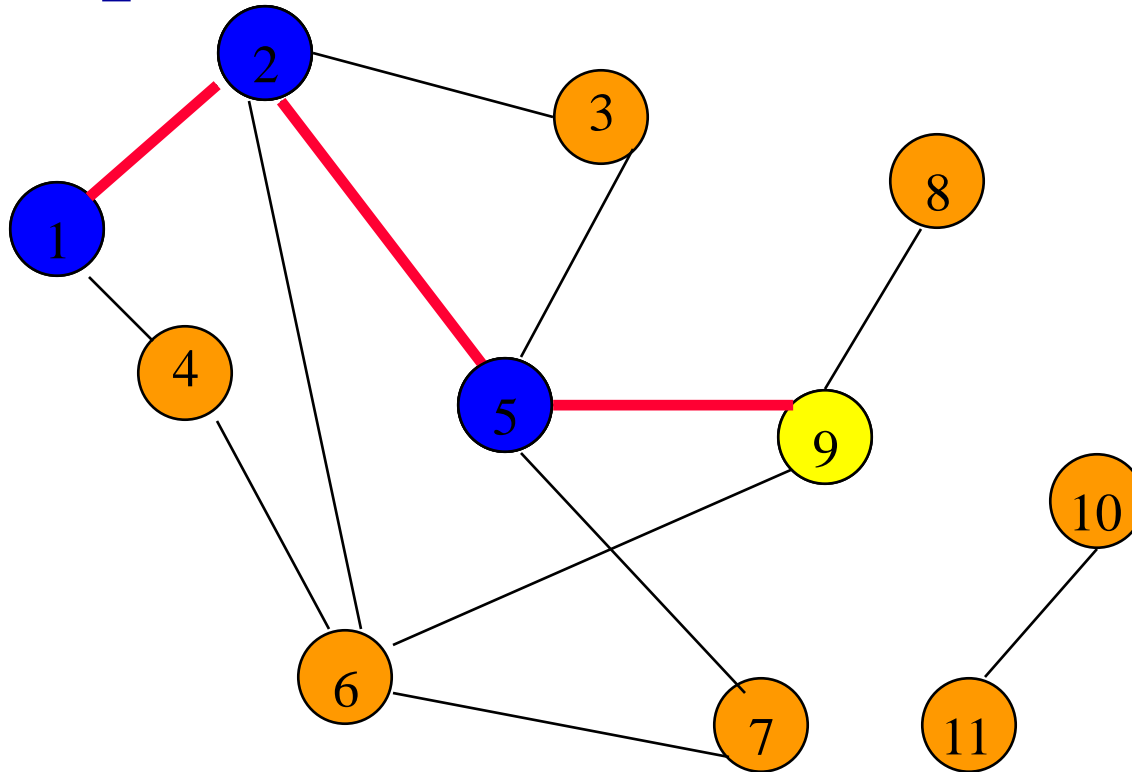
# Depth-First Search Example



Label vertex **2** and do a depth first search from either **3**, **5**, or **6**.

Suppose that vertex **5** is selected.

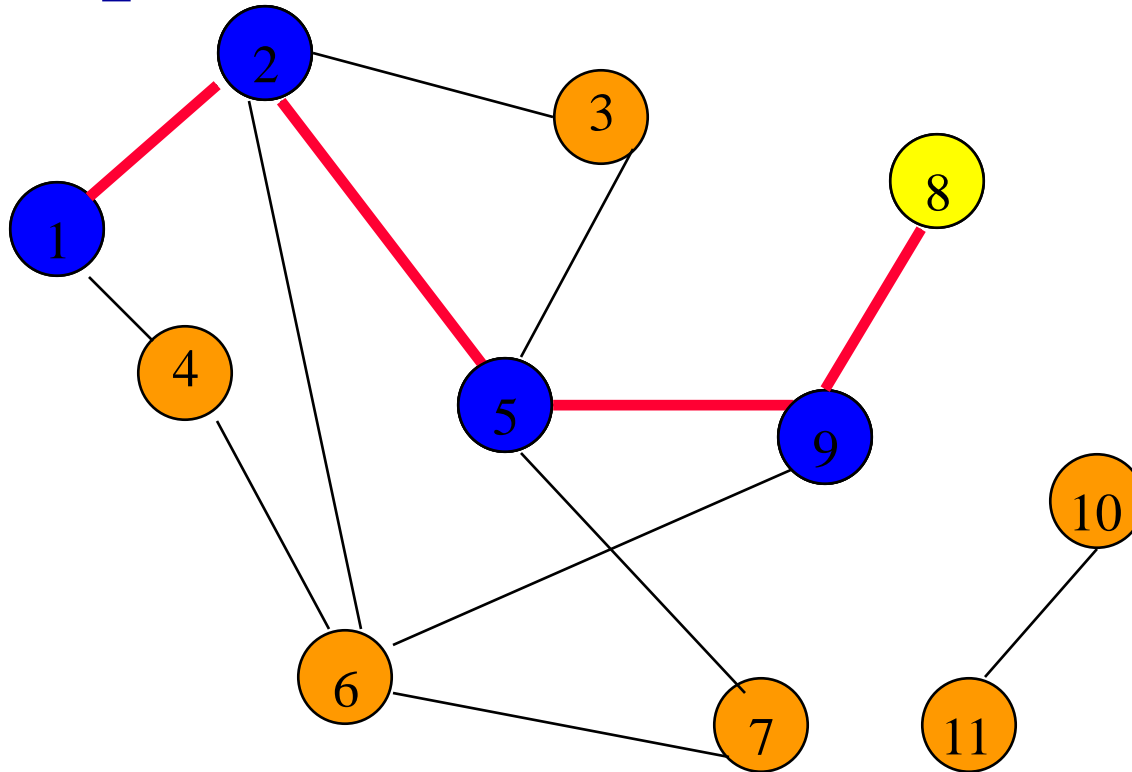
# Depth-First Search Example



Label vertex **5** and do a depth first search from either **3**, **7**, or **9**.

Suppose that vertex **9** is selected.

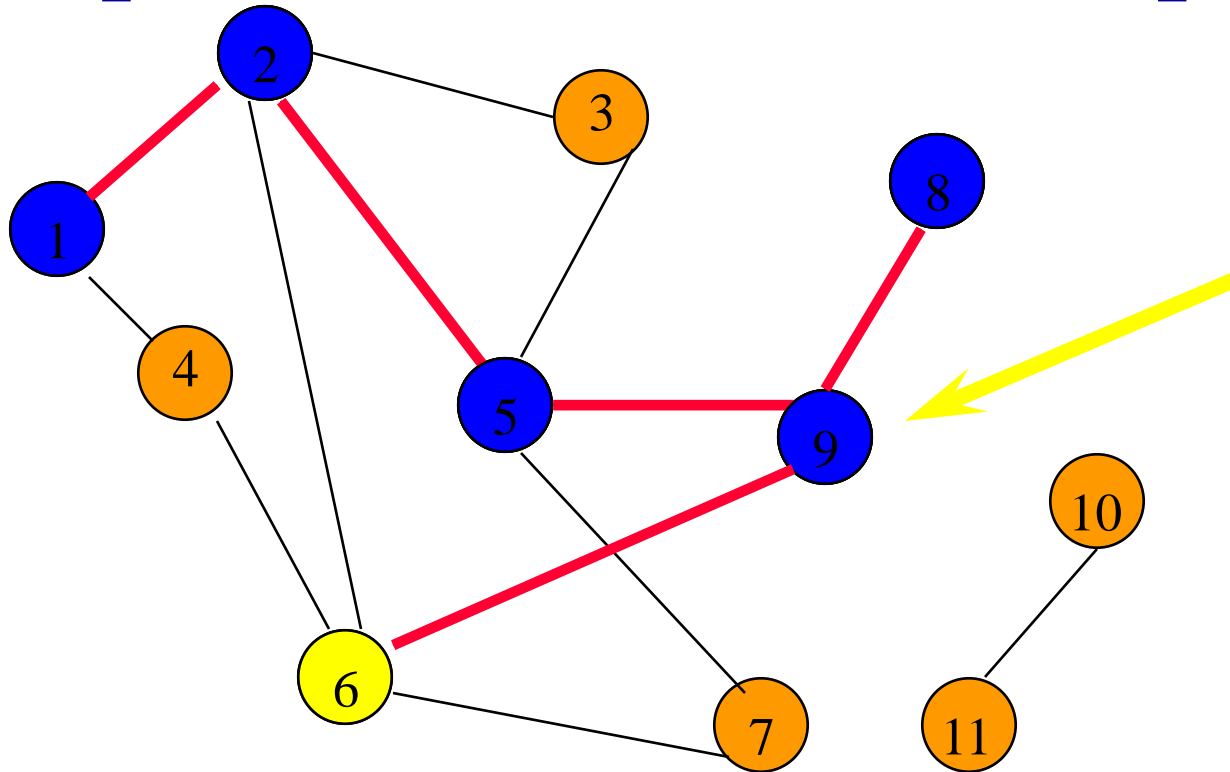
# Depth-First Search Example



Label vertex 9 and do a depth first search from either 6 or 8.

Suppose that vertex 8 is selected.

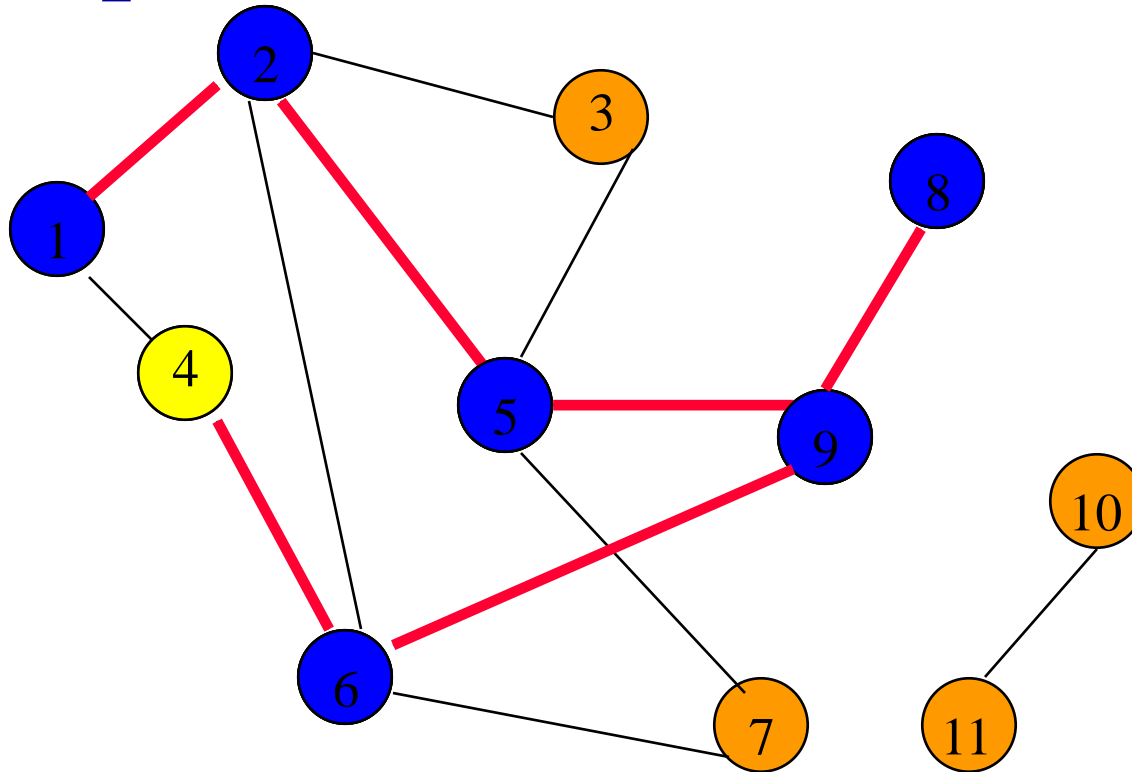
# Depth-First Search Example



Label vertex 8 and return to vertex 9.

From vertex 9 do a  $\text{dfs}(6)$ .

# Depth-First Search Example

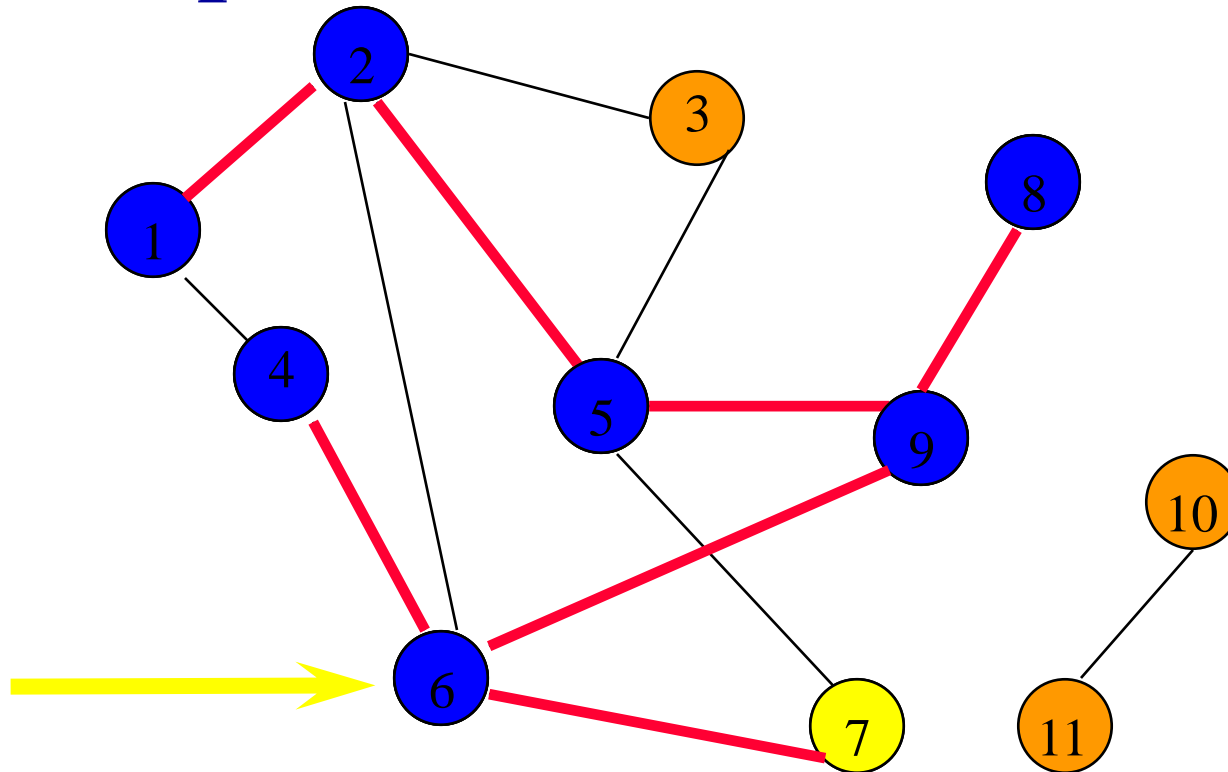


Label vertex **6** and do a depth first search from either **4** or **7**.

Suppose that vertex **4** is selected.



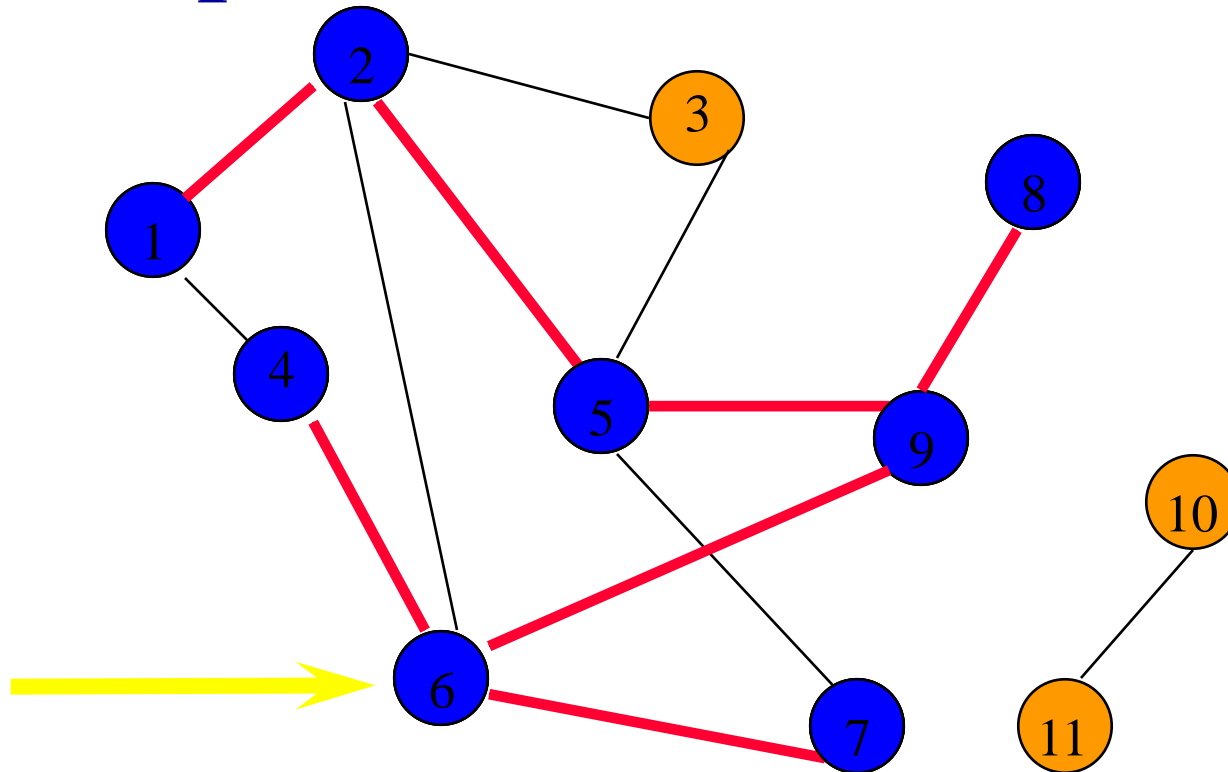
# Depth-First Search Example



Label vertex 4 and return to 6.

From vertex 6 do a  $\text{dfs}(7)$ .

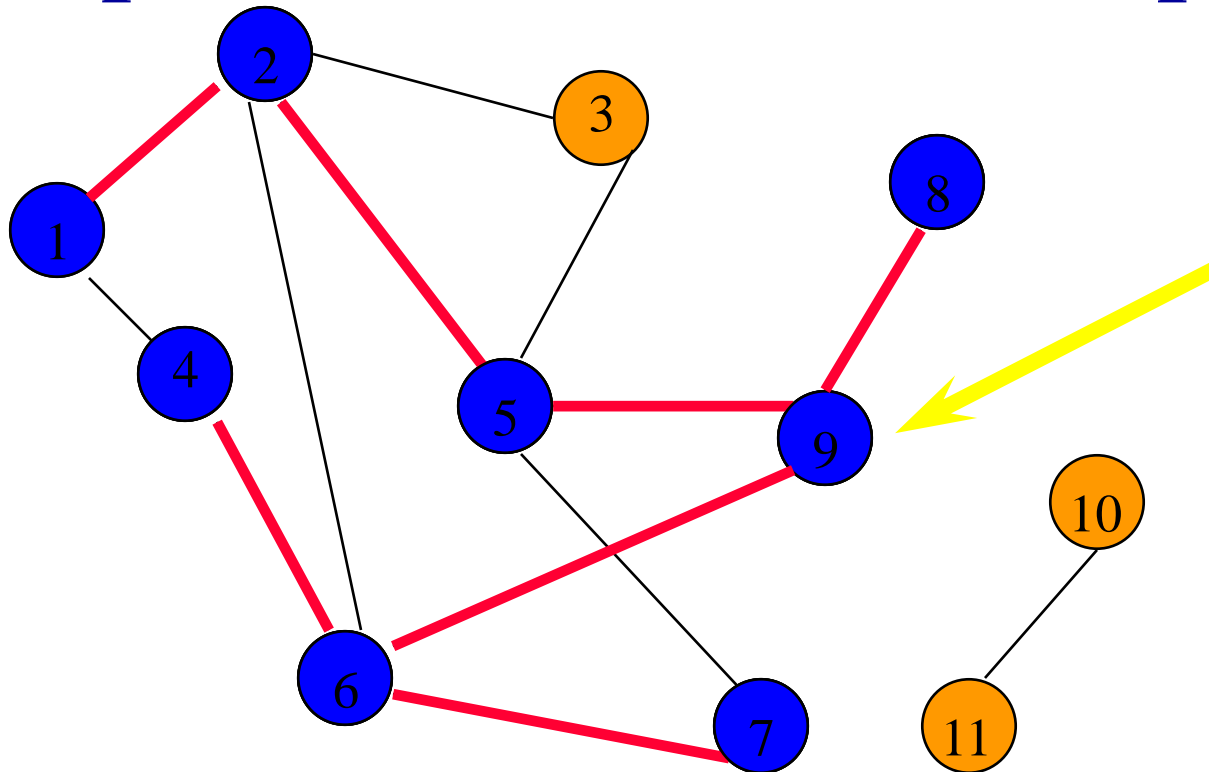
# Depth-First Search Example



Label vertex **7** and return to **6**.

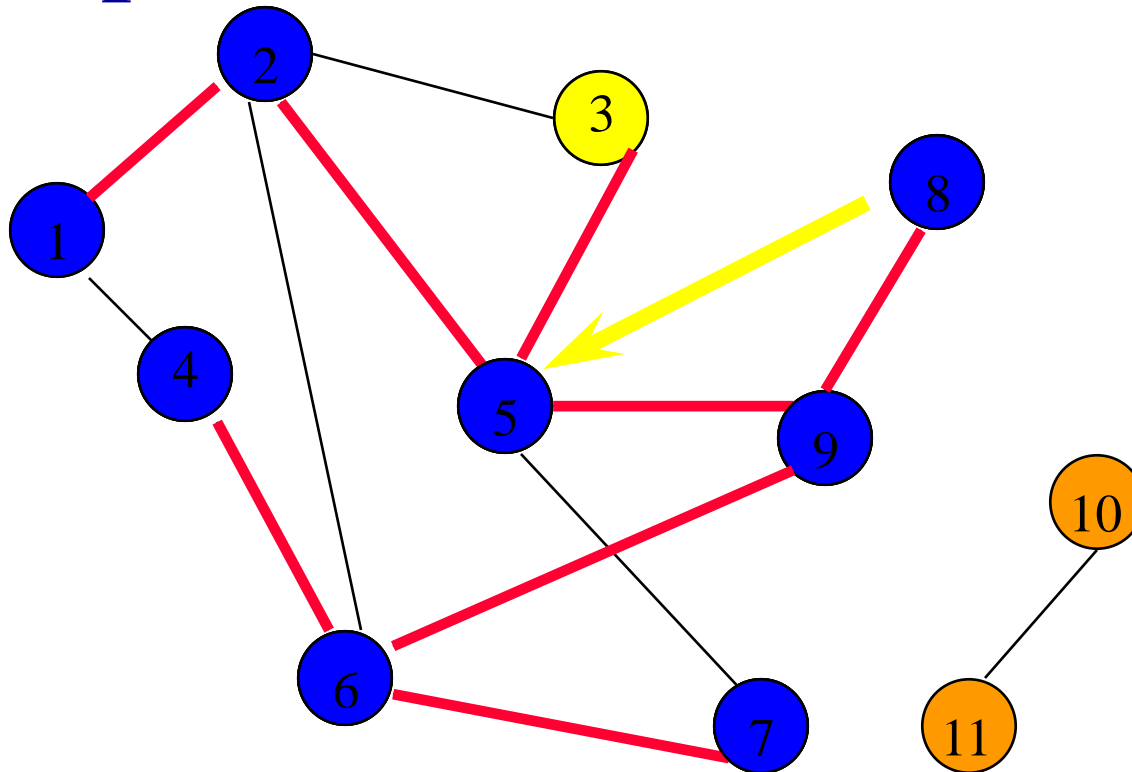
Return to **9**.

# Depth-First Search Example



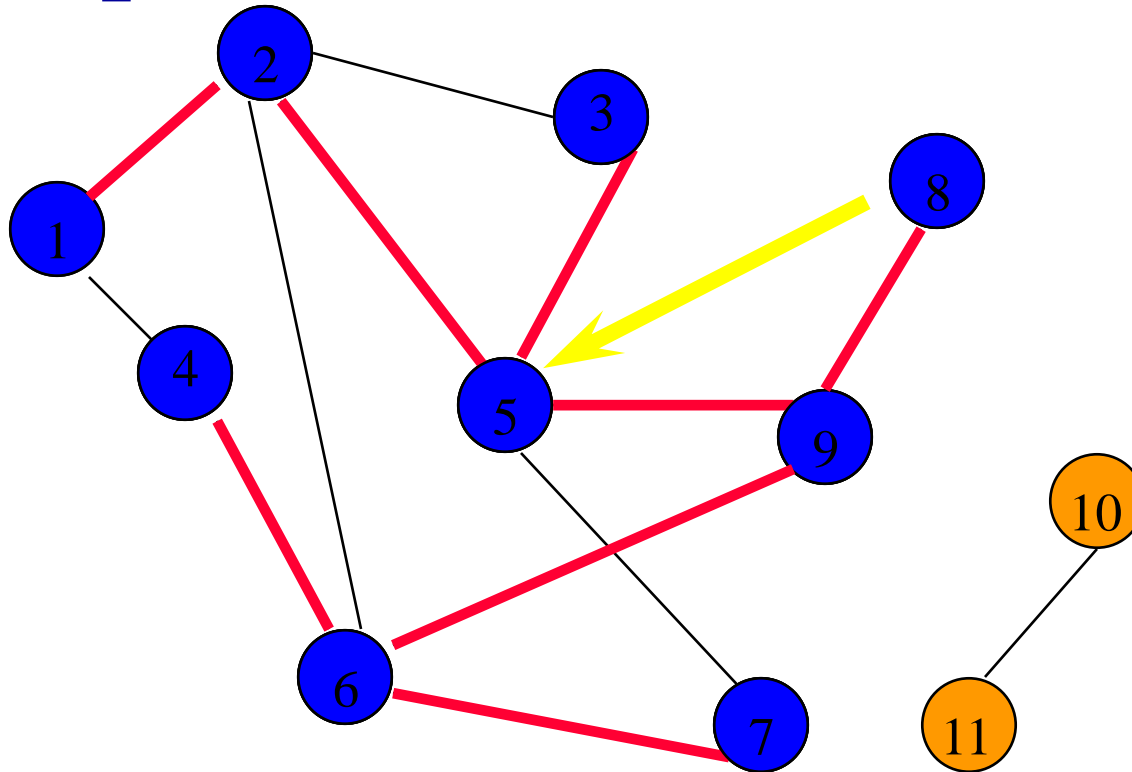
Return to 5.

# Depth-First Search Example



Do a  $\text{dfs}(3)$ .

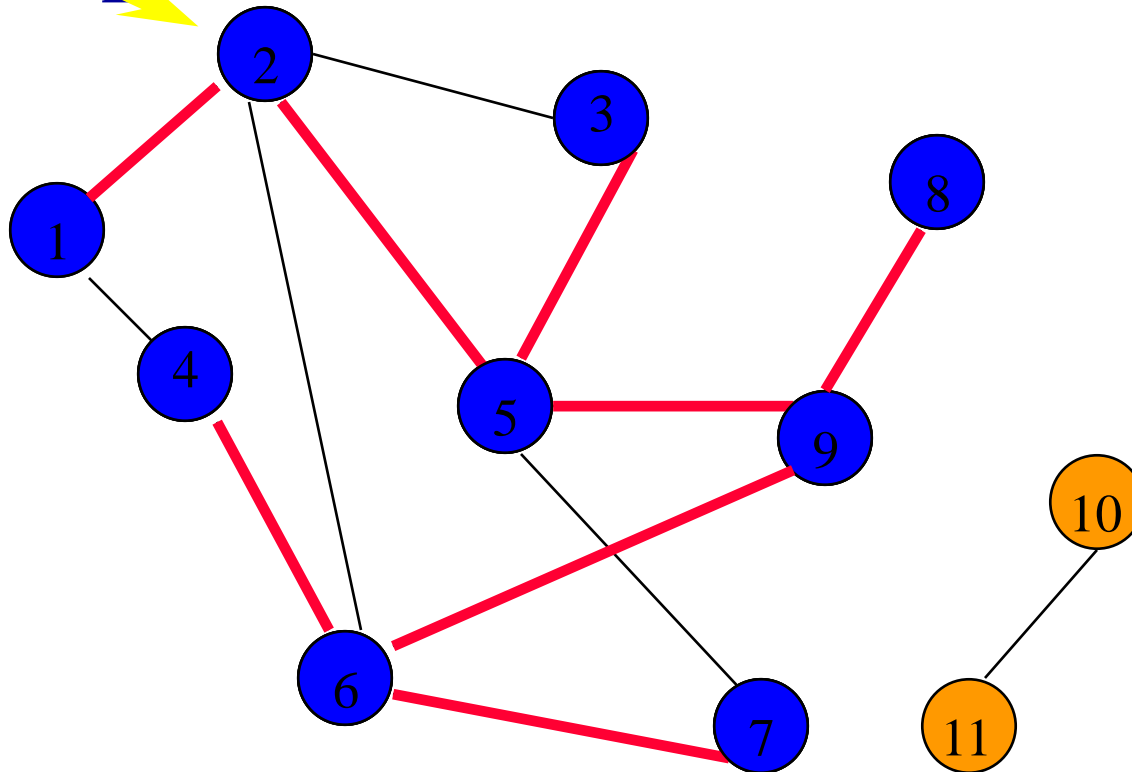
# Depth-First Search Example



Label **3** and return to **5**.

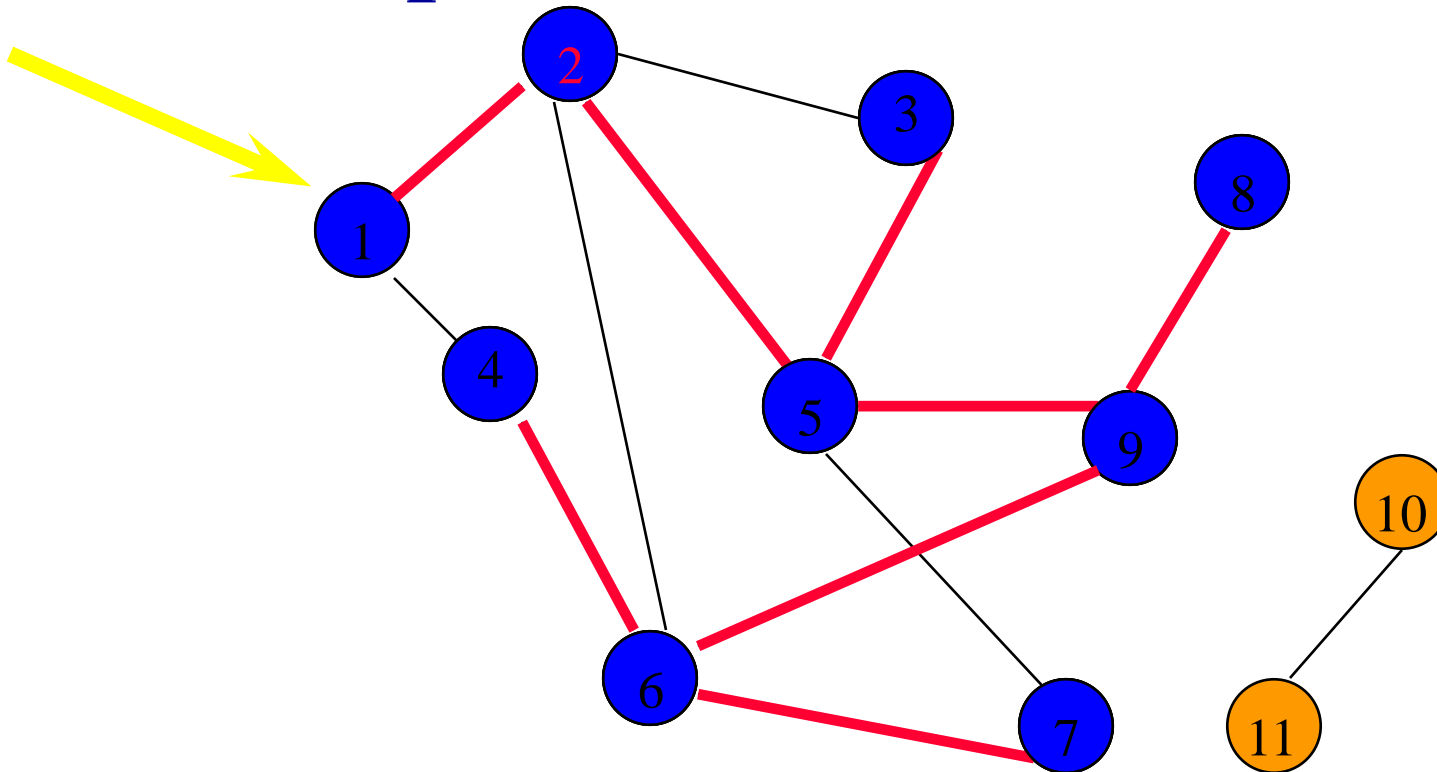
Return to **2**.

# Depth-First Search Example



Return to **1**.

# Depth-First Search Example



Return to invoking method.

# Depth-First Search Properties

- Same complexity as BFS.
- Same properties with respect to path finding, connected components, and spanning trees.
- Edges used to reach unlabeled vertices define a depth-first spanning tree when the graph is connected.
- There are problems for which bfs is better than dfs and vice versa.