

Divide & Conquer

TDT4121 Assignment Lecture 28/09



Part 1

- Understand the divide and conquer algorithm design method
- Understand mergesort
- Understand how ideas from mergesort are used to solve the Counting Inversions problem
- Understand the Finding the Closest Pair of Points problem



Understand the divide and conquer algorithm design method – Three steps

- 1. Break the input into several parts
- 2. Solve the problem in each part recursively
- 3. Combine these solutions to an overall solution



$$5^*2+3/6$$



$$5*2+3/6$$

→ 1. $5*2$ AND $3/6$

→ 2. 10 AND 0.5

→ 3. 10.5



Understand Mergesort





Merge Sort in 3

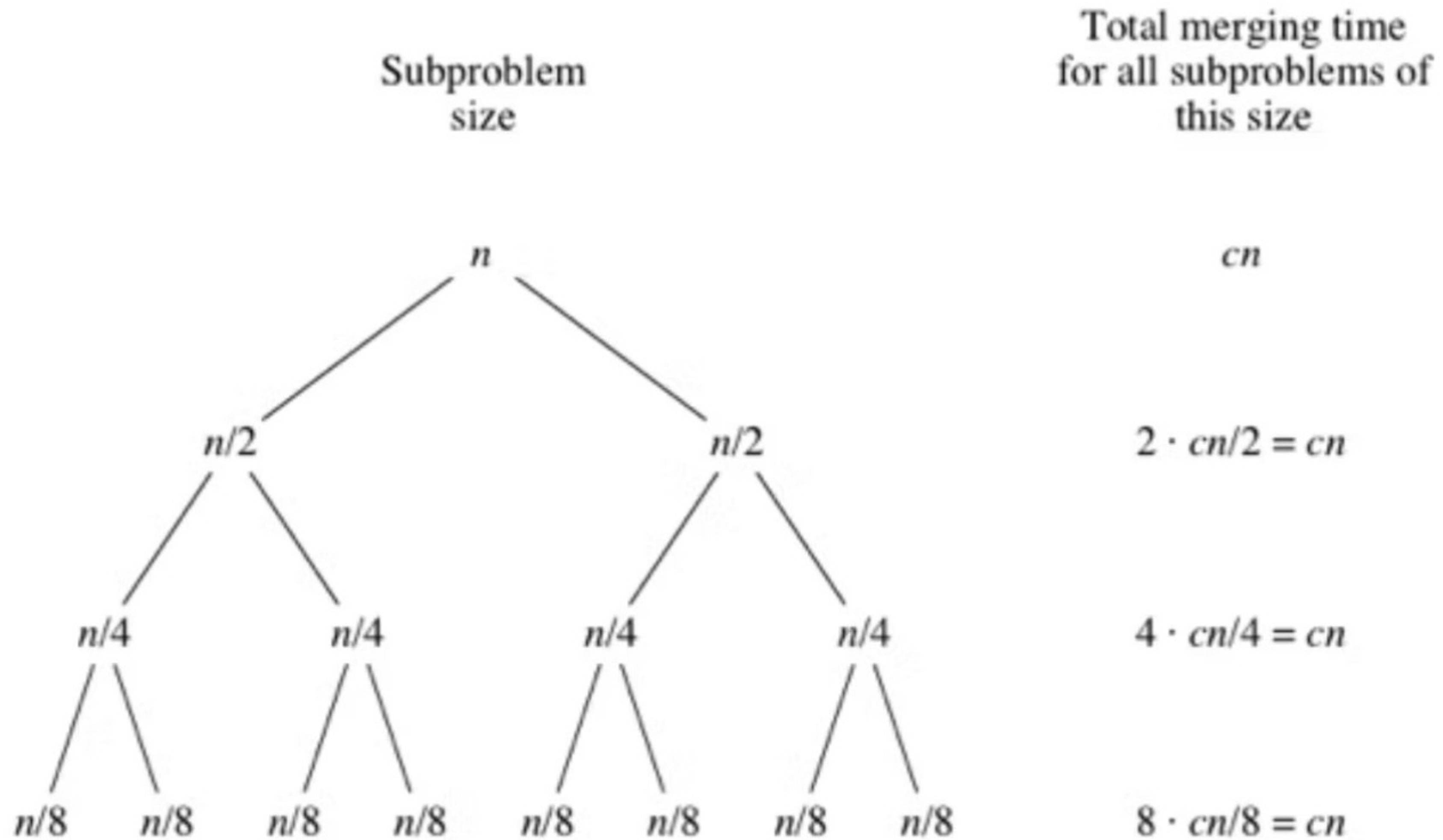
Watch on  YouTube



Runtime

- 1. The divide step takes constant time, regardless of the subarray size: $O(1)$
- 2. The conquer step, where we recursively sort two subarrays of approximately $n/2$ elements each, takes some amount of time
- 3. Combining these three takes $O(n)$ time, by merging together n elements.

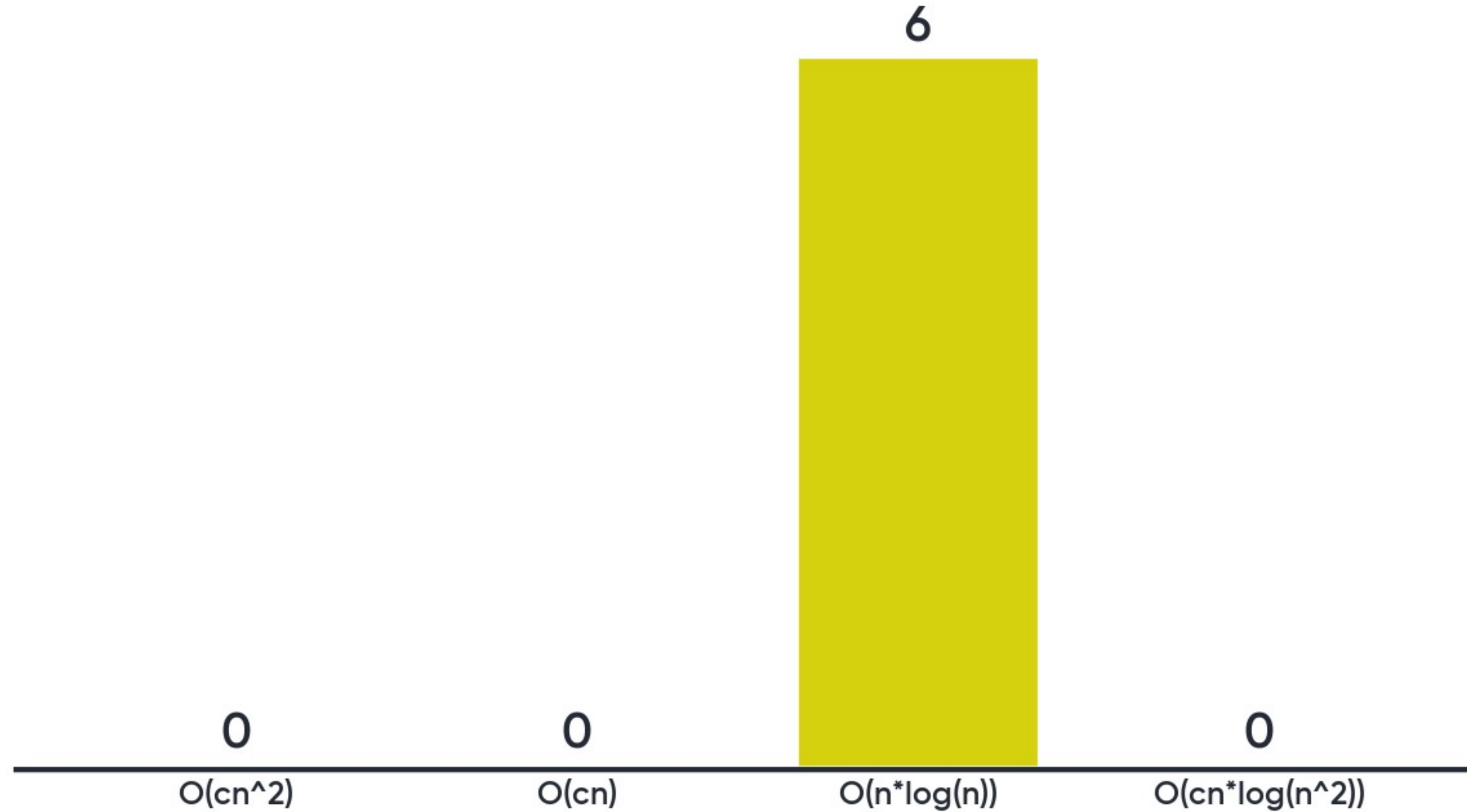




Hint: $cn \cdot \text{height of tree} = \text{Runtime}$



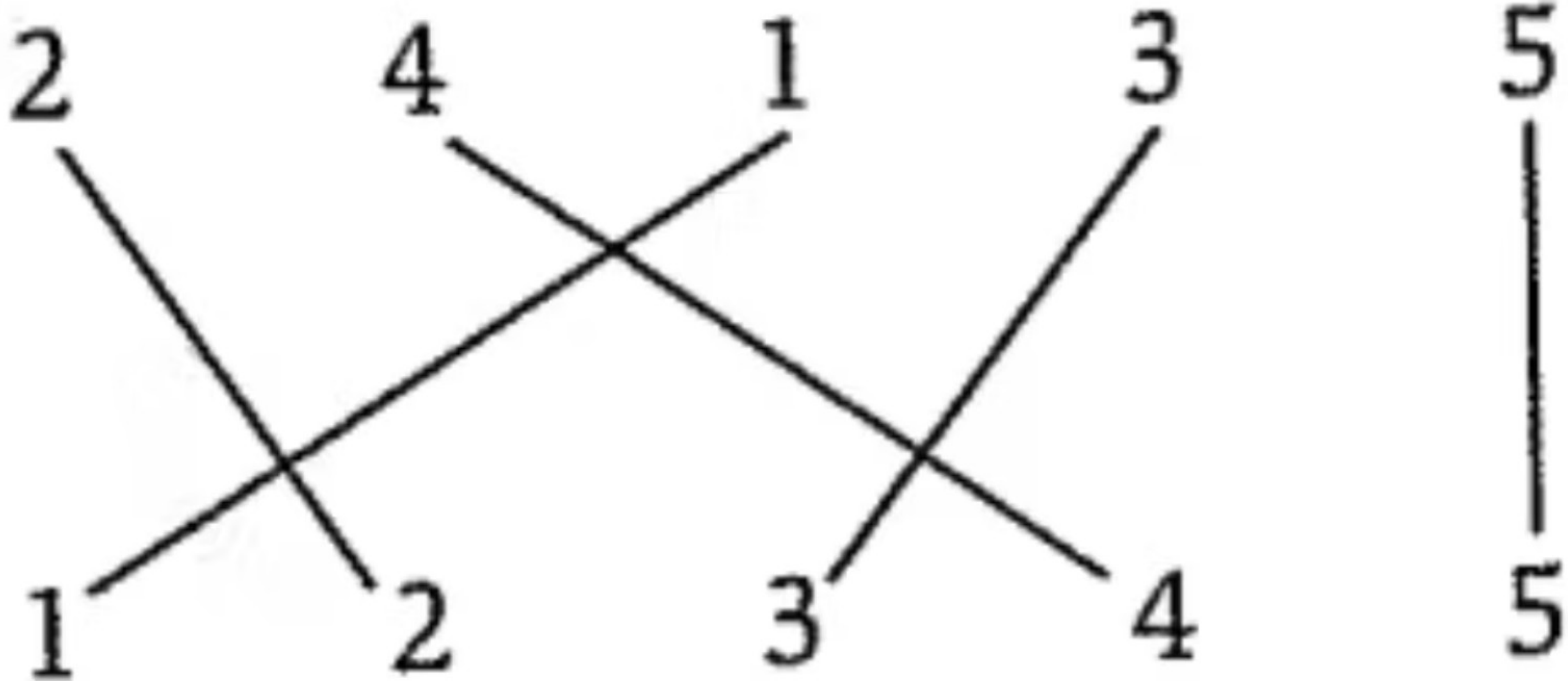
What is the runtime of Mergesort?



Understand how ideas from mergesort are used to solve the Counting Inversions problem

- 1. The Counting Inversions problem
- 2. Using Mergesort to solve it
- 3. Runtime





inversions are pairs of numbers, in a disordered list, where the larger of the two numbers is to the left of smaller number.



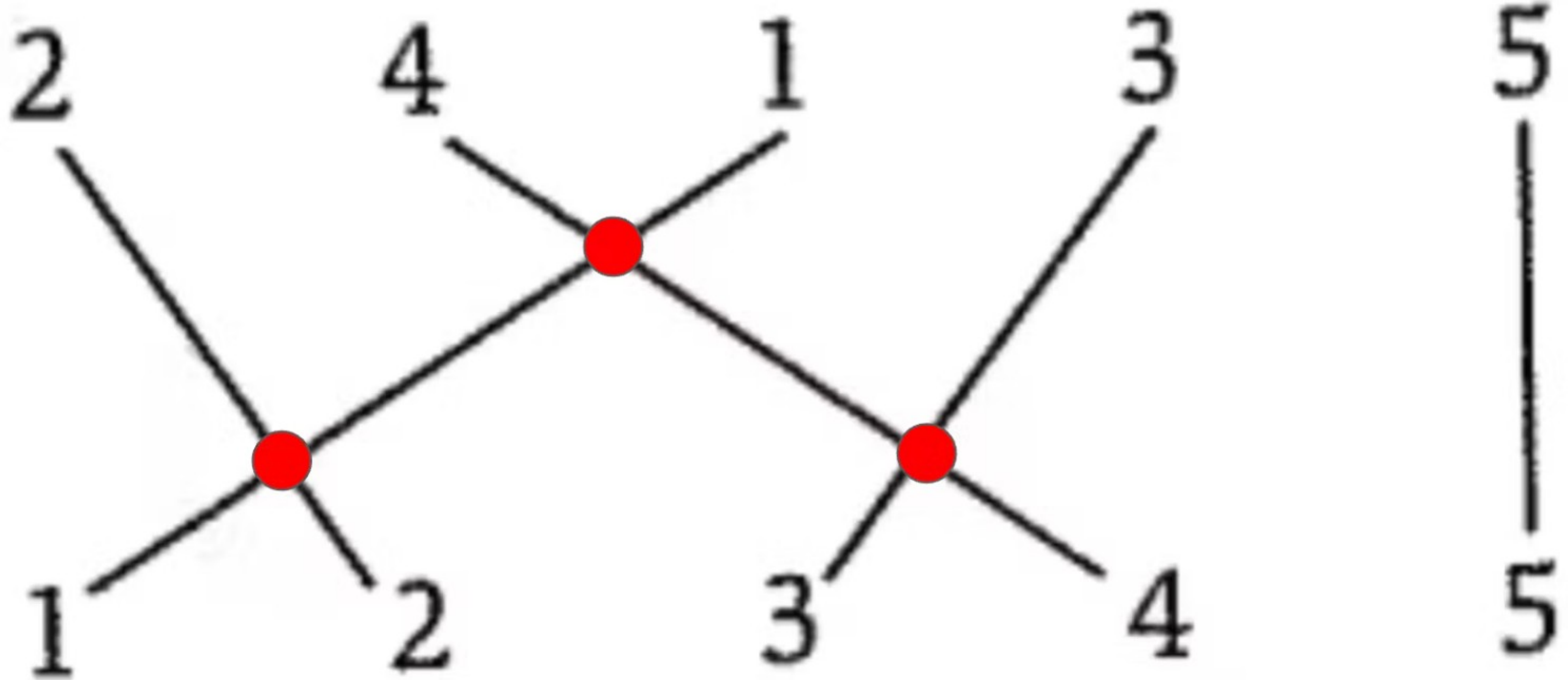
How many inversions?

0
1

0
2

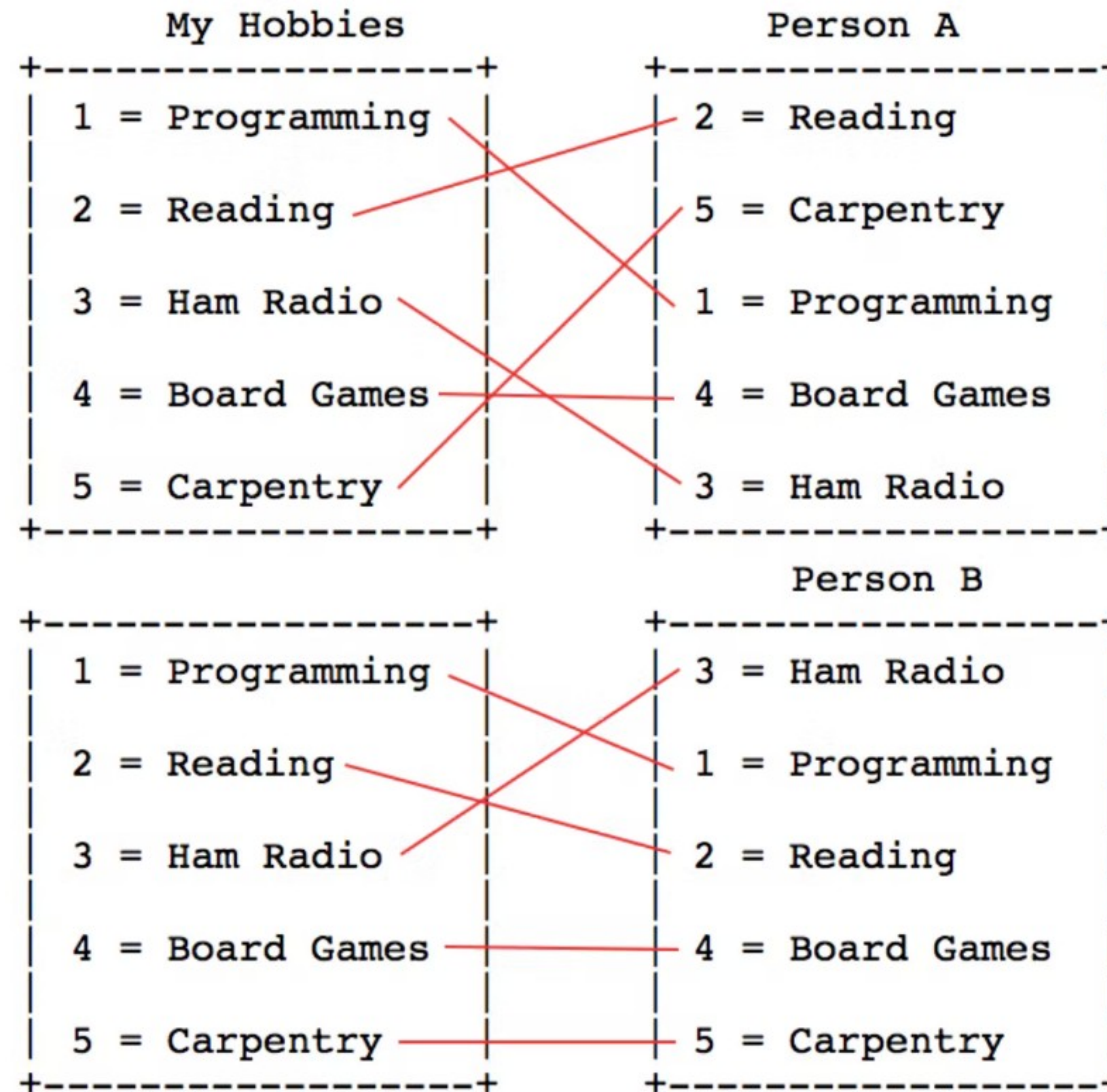
5
3

0
4



Trick: Count crosses





Task: Who should I hang out with?



Who should I hang out with?

0
Person A


Person B

Sort-and-Count(L)

If the list has one element then
there are no inversions

Else

Divide the list into two halves:

A contains the first $\lfloor n/2 \rfloor$ elements

B contains the remaining $\lfloor n/2 \rfloor$ elements

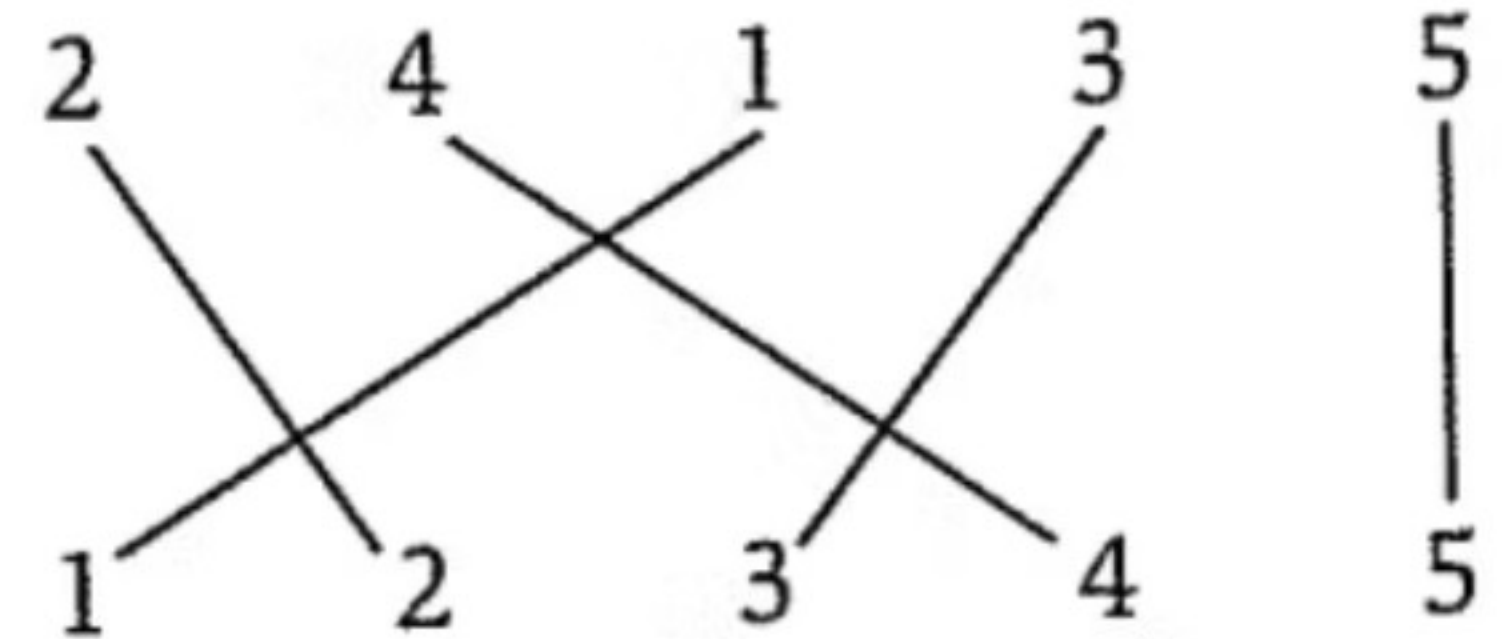
$(r_A, A) = \text{Sort-and-Count}(A)$

$(r_B, B) = \text{Sort-and-Count}(B)$

$(r, L) = \text{Merge-and-Count}(A, B)$

Endif

Return $r = r_A + r_B + r$, and the sorted list L

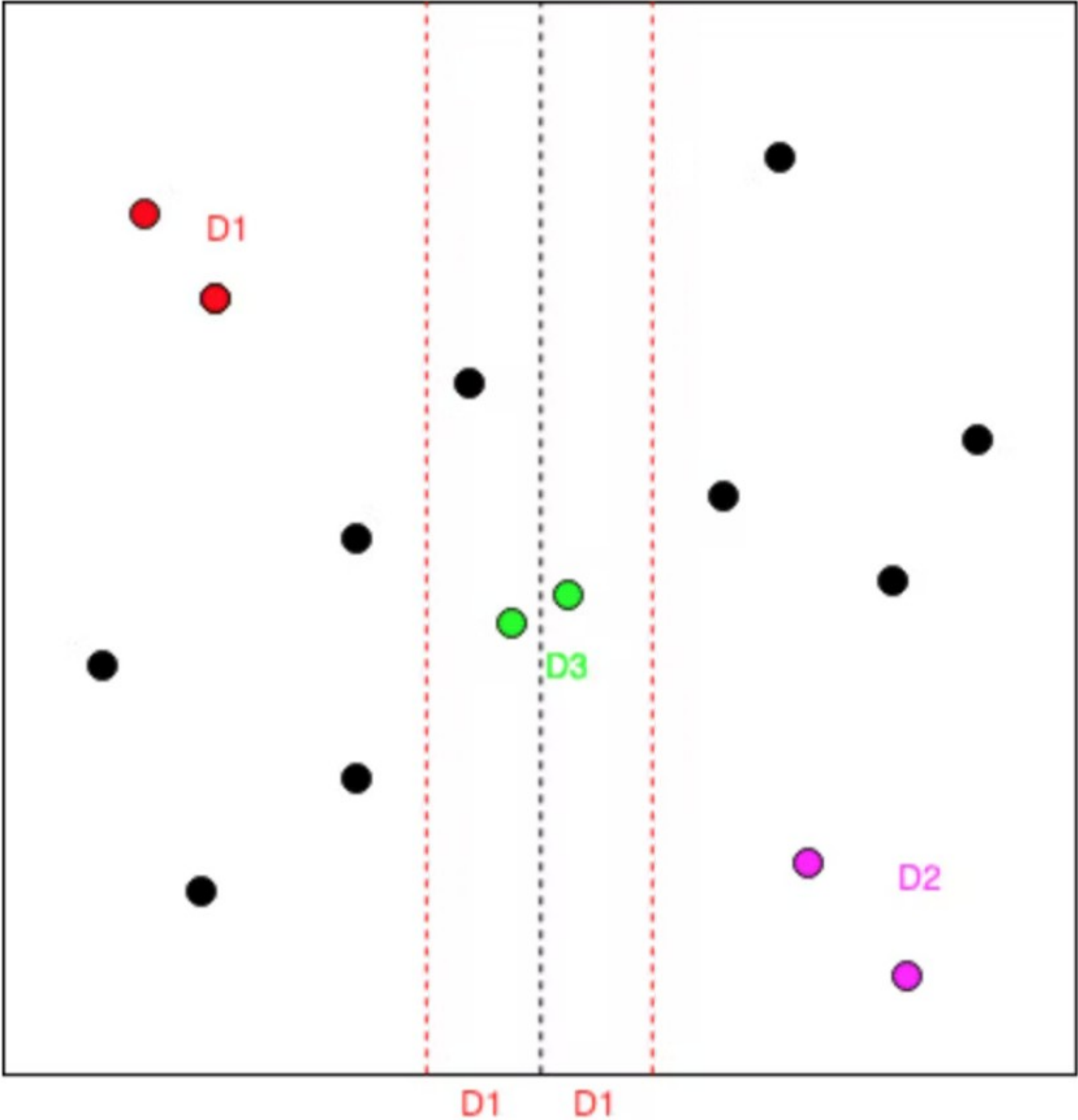


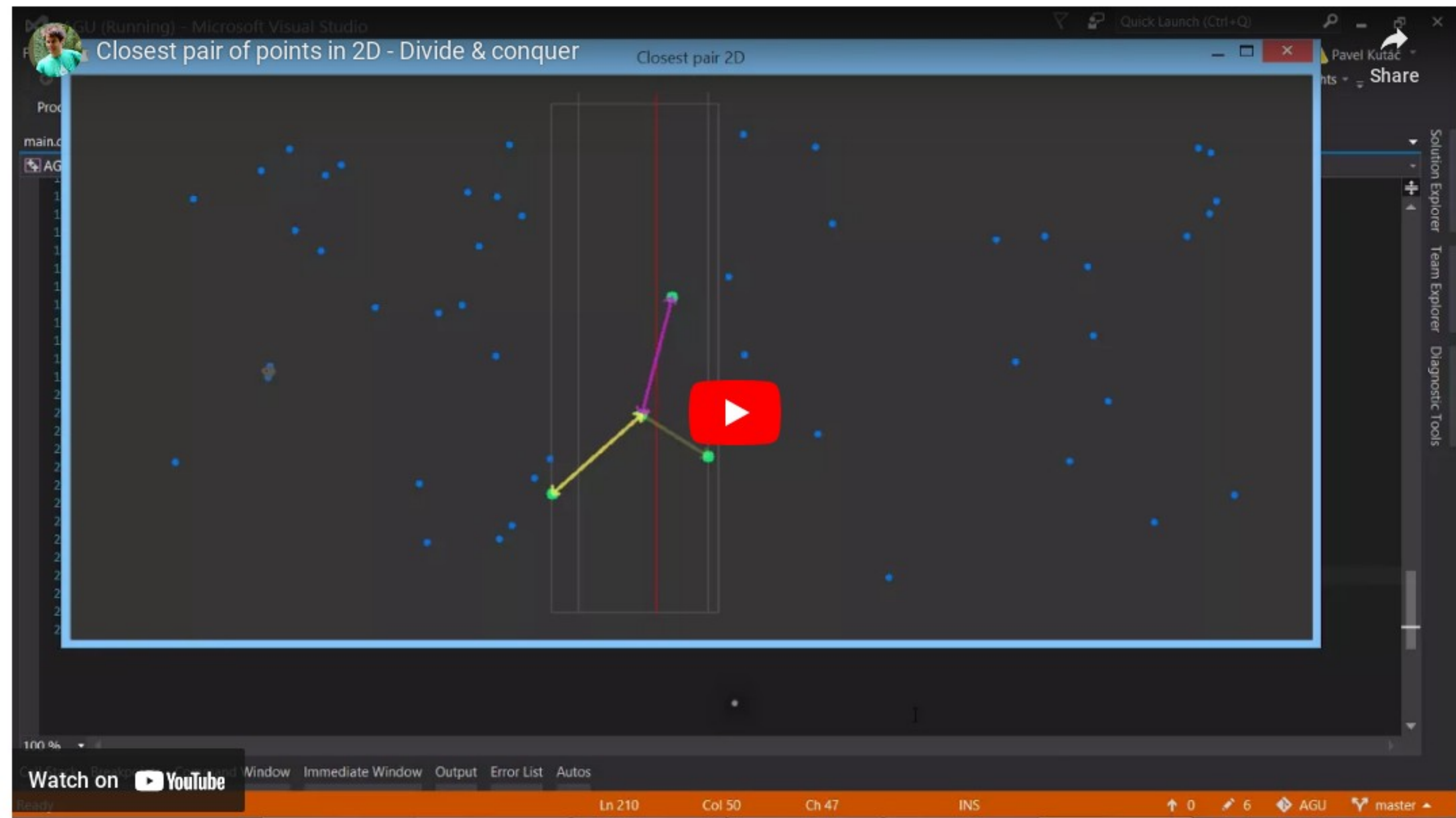
Sort-and-Count algorithm correctly sorts the input list and counts the number of inversions; it runs in $O(n \log n)$ time for a list with n elements.



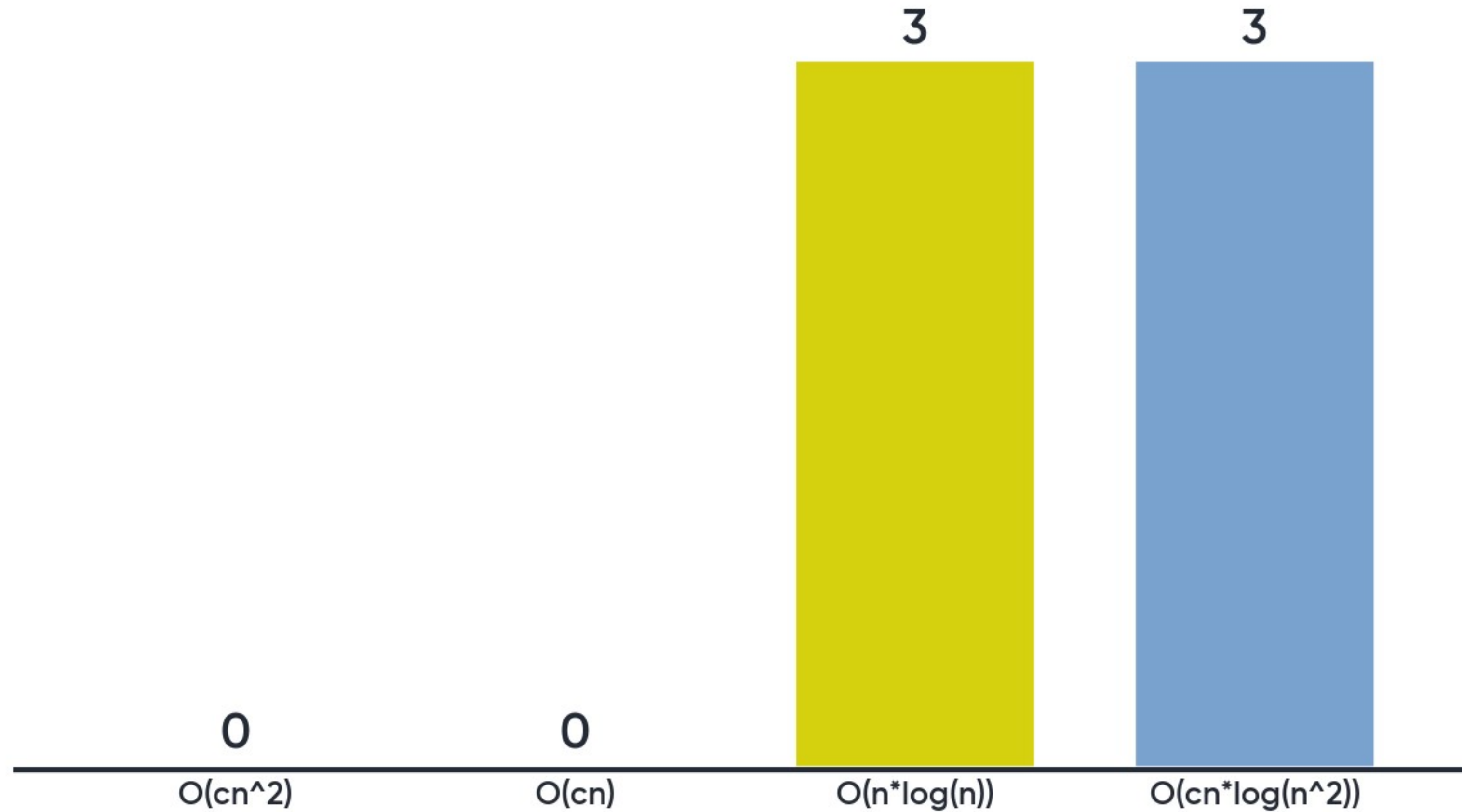
Understand the Finding the Closest Pair of Points problem







What is the runtime of CPOP?



Part 2

- Understand the Integer Multiplication problem
- Solve a recurrence by unrolling
- Solve a recurrence by induction
- Understand the Strassen algorithm



The Integer Multiplication problem

→ Grade school: $O(n^2)$ 🙄

→ Karatsuba-Ofman: $O(n^{1.585})$ 👍



Recursive-Multiply(x, y):

Write $x = x_1 \cdot 2^{n/2} + x_0$

$y = y_1 \cdot 2^{n/2} + y_0$

Compute $x_1 + x_0$ and $y_1 + y_0$

$p = \text{Recursive-Multiply}(x_1 + x_0, y_1 + y_0)$

$x_1 y_1 = \text{Recursive-Multiply}(x_1, y_1)$

$x_0 y_0 = \text{Recursive-Multiply}(x_0, y_0)$

Return $x_1 y_1 \cdot 2^n + (p - x_1 y_1 - x_0 y_0) \cdot 2^{n/2} + x_0 y_0$

Karatsuba-Ofman Integer Multiplication



Exercises: Karatsuba

1. Solve $36 \cdot 28$

2. Solve $101_2 \cdot 100_2$



Using Karatsuba, what is $11_2 \cdot 11_2$?

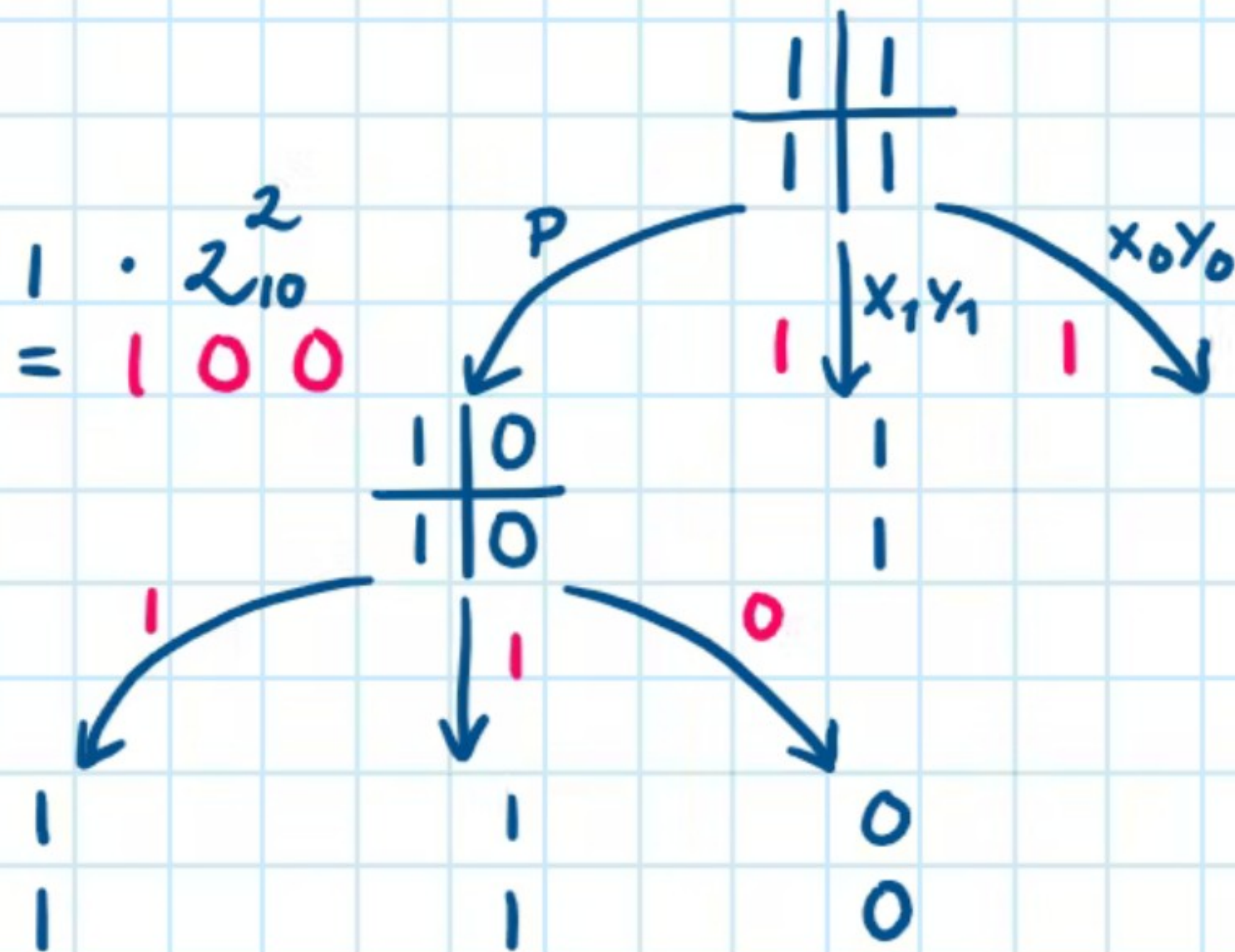
0
1010


1
1001

0
1111

0
0101

$$1 \cdot 2_{10}^2 + (100 - 1 - 1) \cdot 2_{10} + 1 \\ = 100 + 100 + 1 = \underline{\underline{1001}}$$



What recurrence describes the Karatsuba algorithm?

0
 $T(n) = 4T(n/2) + cn$

2
 $T(n) = 3T(n/2) + cn$

1
 $T(n) = 3T(n/3) + cn$

0
 $T(n) = 4T(n/3) + cn$

Exercises: Recurrences

1. Solve the Karatsuba recurrence by unrolling
2. Confirm your answer by substitution



Solve a recurrence by unrolling

1. Analyze the first few levels
2. Identify a pattern
3. Sum over all levels of recursion



Solve a recurrence by induction

1. Inductive hypothesis
2. Show inductive step



Exercises: Complexity

1. Plot the number of operations of Karatsuba for $1 < n < 900$
2. Confirm by plot that the complexity is $O(n^2)$
3. Confirm by plot that the complexity is $\Theta(n^{1.585})$



The Strassen algorithm

→ Matrix Multiplication in $O(n^{2.81})$



Exercises: Matrix Multiplication

- Implement a naive algorithm and find the complexity
- Implement a D&Q algorithm and find the recurrence
- Implement Strassen and find the recurrence



What is the complexity of the naive matrix multiplication algorithm?

$$O(n^2)$$

$$O(n \cdot \log(n))$$

$$O(n^{\log(n)})$$


$$O(n^3)$$

What recurrence describes the D&Q matrix multiplication algorithm?

$$T(n) = 8T(n/2) + cn$$

$$T(n) = 4T(n/2) + cn^2$$



$$T(n) = 8T(n/2) + cn^2$$

$$T(n) = 4T(n/4) + cn$$

What recurrence describes the Strassen matrix multiplication algorithm?



1

$$T(n) = 7T(n/2) + cn^2$$

0

$$T(n) = 8T(n/2) + cn$$

0

$$T(n) = 8T(n/2) + cn^2$$

0

$$T(n) = 7T(n/2) + cn$$

2.12. How many lines, as a function of n (in $\Theta(\cdot)$ form), does the following program print? Write a recurrence and solve it. You may assume n is a power of 2.

```
function f(n)
    if n > 1:
        print_line('still going')
        f(n/2)
        f(n/2)
```



What recurrence describes the number of lines that will be printed?

$$T(n) = T(2n), T(1) = 1$$



$$T(n) = 2T(n/2) + 1, T(1) = 0$$

$$T(n) = 2T(n), T(0) = 1$$

$$T(n) = 2T(n/4) - 1, T(0) = 1$$

How many lines will be printed?



$\Theta(\log(n))$

0

$\Theta(n \cdot \log(n))$

0

$\Theta(n)$

0

$\Theta(n^2)$

Ask me anything

0 questions
0 upvotes